Efficient In-Memory Inverted Indexes Theory and Practice



Joel Mackenzie The University of Queensland



Sean MacAvaney University of Glasgow



Antonio Mallia Pinecone



Michal Siedlaczek MongoDB, Inc.







Who are we?



Joel Mackenzie
The University of
Queensland



Sean MacAvaney University of Glasgow



Antonio Mallia Pinecone



Michal Siedlaczek MongoDB, Inc.

Early Career Academics/Practitioners interested in efficient and effective IR systems. All very keen on empirical, "hands-on" research and development.

https://jmmackenzie.io/ https://macavaney.us/ https://www.antoniomallia.it/ https://siedlaczek.me/

Who are you?

- Newcomers with no assumed knowledge of efficient inverted index-based architectures and query processing.
- **Experienced Researchers** who are keen to sharpen their skills and improve their understanding of efficient indexing and retrieval.
- Everyone in between!

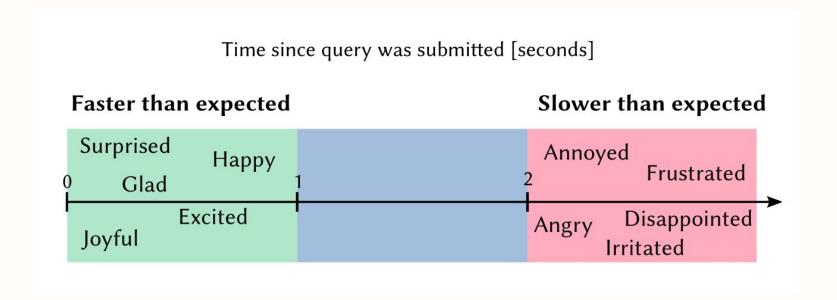
Why should you care?

- Despite the best efforts of the IR community, inverted indexes just will not die!
 - Sometimes you really do need to find documents that contain a specific set of terms;
 - Traditional ranking models like BM25 continue to be a strong baseline on unseen data;
 - Inverted indexes tend to scale extremely well as collections grow large.

Why should you care?

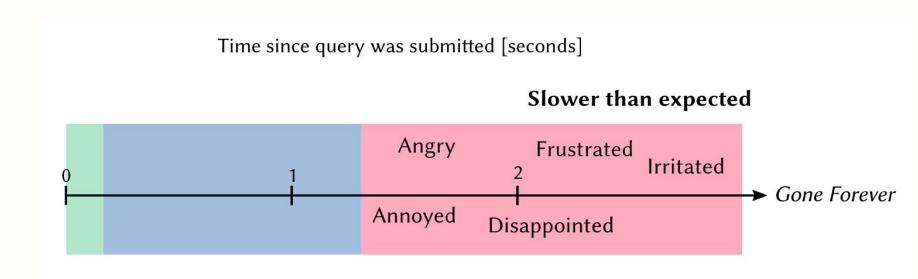
- Next generation ranking methods such as learned sparse retrieval are still making use of inverted indexes.
- Possessing a strong theoretical understanding of inverted indexes, including how they can be engineered to be efficient in practice, is still very relevant in 2025!
- This tutorial will cover both the theory and practice necessary to understand, experiment with, and contribute to the future of inverted indexes.

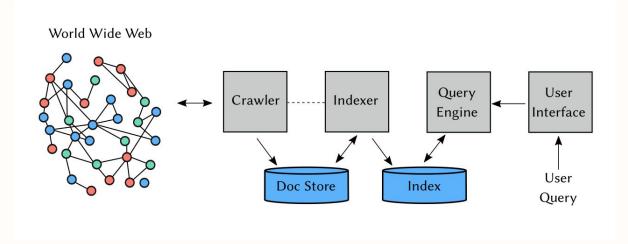
"returning good results quickly is better than returning the best results slowly" – Dean and Barosso, CACM, 2013.

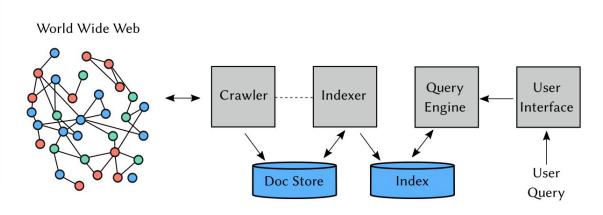


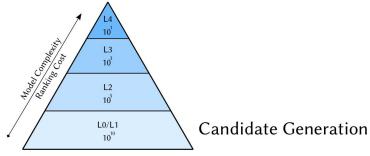
"returning good results quickly is better than returning the best results slowly"

- Dean and Barosso, CACM, 2013.











https://www.google.com/about/datacenters/gallery

Today's Plan

Session 1: (Basic) Indexing and Retrieval 09.00 - 10.30

Morning Tea 10.30 - 11.00

Session 2: Learned Sparse Retrieval 11.00 - 12.00

Session 3: New Directions 12.00 - 12.30

Intended Learning Outcomes

ILO 1: Theoretical Understanding of Inverted Indexes

ILO2: Fast Top-k Retrieval with Dynamic Pruning

ILO3: Current Trends and Research Directions

ILO4: Experimenting with PISA

ILO5: Integrating PISA into Modern Applications

IR Experimentation













History and Origins of PISA

The **PISA engine** started off as the *Data Structures for Inverted Indexes (ds2i)* project in 2014.

It formed the basis of "Partitioned Elias Fano Indexes" by Ottaviano and Venturini, which won the **SIGIR 2014** best paper award, as well as "Optimal Space-Time Tradeoffs for Inverted Indexes" by Ottaviano, Tonellotto, and Venturini, **WSDM 2015.**

In 2017, "Faster BlockMax WAND with variable-sized blocks" was published at **SIGIR** by Mallia, Ottaviano, Porciani, Tonellotto, and Venturini. At this point, **PISA was forked from ds2i.**

The PISA Engine

Q PISA

An efficient, extensible, modern search engine.

- Written in C++17
- In-memory retrieval
- Low-level optimization out-of-the-box: CPU intrinsics, branch prediction hinting, ...
- Extensible: Plug and play parsing, stemming, compression, query processing
- Indexing, parsing, sharding capabilities
- Free, open-source permissive license

Where PISA Shines

Ridiculously fast top-k query processing



Extensible experimentation with easy access to state-of-art methods

Small but active group of maintainers

Interfaces well with other experimental IR systems

Where PISA Pales

Q PISA

Primary focus is on bag-of-words, top-k retrieval

No support for positional indexing, fields, ...

Not so user friendly due to high complexity of the codebase

> See: https://github.com/terrierteam/pyterrier_pisa for a higher-level Python API



PISA's role in the IR Ecosystem



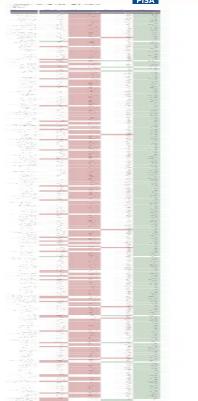
PISA, because it happily trades anything in favor of faster search, so
it sets a good north star in terms of search performance.

https://www.reddit.com/r/rust/comments/nu1jc7/tantivy_v015_released_now_backed_by_quickwit_inc/https://jpountz.github.io/2025/05/12/analysis-of-Search-Benchmark-the-Game.html

Search Benchmark, the Game Collection type TOP_1000 Type of Query union ▶ tantivy-0.22 ▶ lucene-10.2.0 ▶ lucene-10.2.0-bp ▶ pisa-0.8.2 Query 3,907 µs 702 µs 1,429 µs 1,104 µs 121 µs griffith observatory +812.4% +480.2 % 1 docs 1 docs 296 µs 733 µs 547 µs 80 µs bowel obstruction +583.8 % 1 docs 1 docs 1 docs

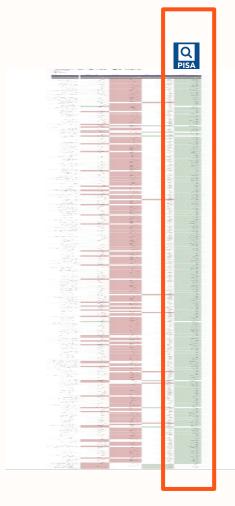






Search Benchmark, the Game Collection type TOP_1000 Type of Query union ▶ tantivy-0.22 ▶ lucene-10.2.0 ▶ lucene-10.2.0-bp ▶ pisa-0.8.2 Query 3,907 µs 702 µs 1,429 µs 1,104 µs 121 µs griffith observatory +812.4% +480.2 % 1 docs 1 docs 296 µs 733 µs 547 µs 80 µs bowel obstruction +583.8 % 1 docs 1 docs 1 docs





Session I: Indexing and Retrieval

Setup

Links and Downloads

For the practical component, we will need to grab some data, and have some instructions ready.

Prerequisite: You have your own machine with Docker installed. I will step through on my own machine if you don't have access.

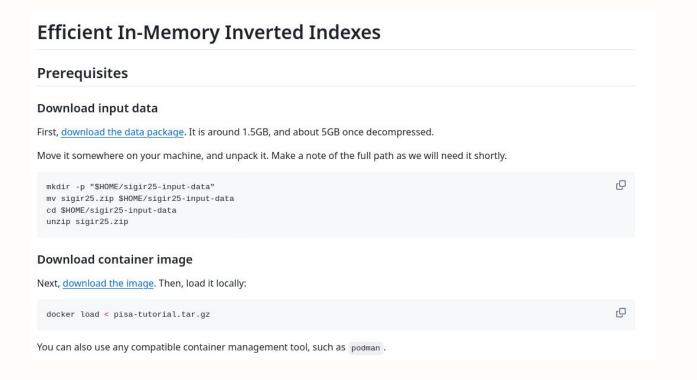
You can also experiment with the tutorial at any time (later).

Tutorial: https://shorturl.at/VExpG

Aka: https://github.com/pisa-engine/pisa/blob/main/test/docker/tutorial/instructions.md

Please at least initiate the data download, and the docker image download.

Links and Downloads



Session I: Indexing and Retrieval

Theory

Document 0

search is cool

Document 1

search is fun

Document 2

search is fun

for everyone

Document 0

search is cool

Document 1

search is fun

Document 2

search is fun for everyone

Lexicon

Postings Lists

Document 0

search is cool

Document 1

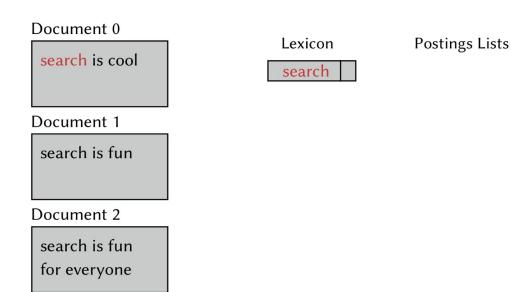
search is fun

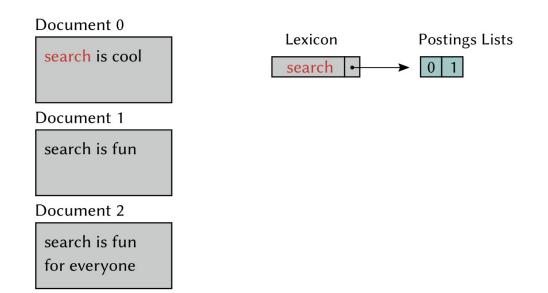
Document 2

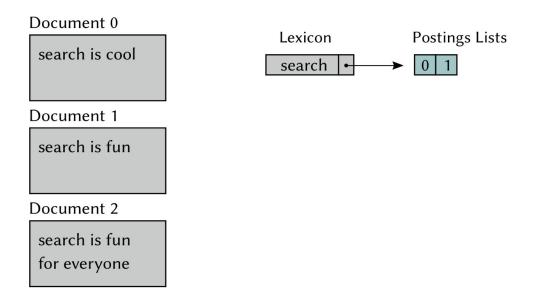
search is fun for everyone

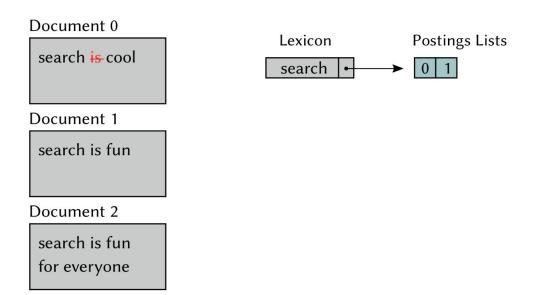
Lexicon

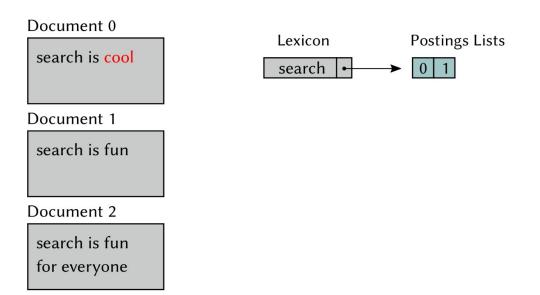
Postings Lists

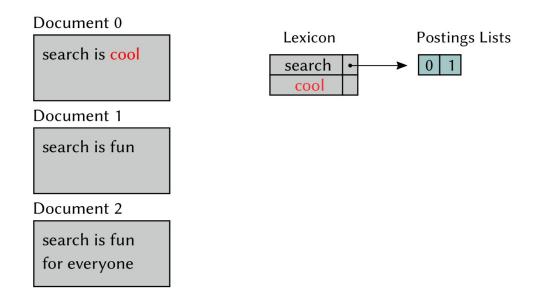




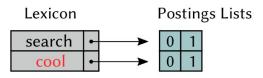


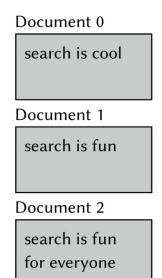


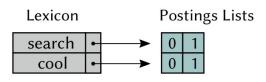














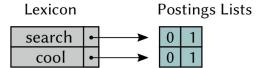
search is cool

Document 1

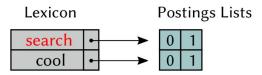
search is fun

Document 2

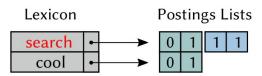
search is fun for everyone

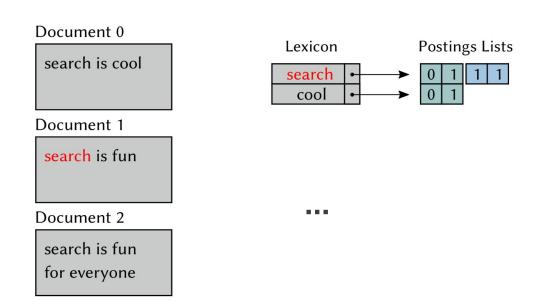


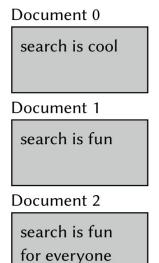


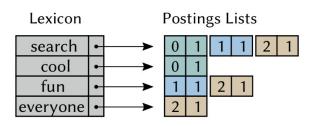


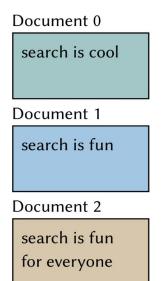


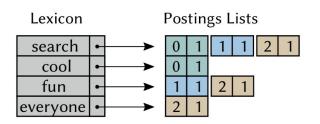


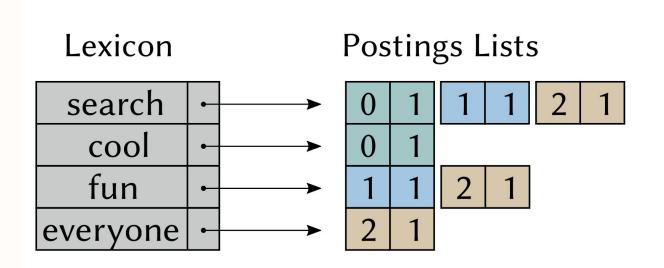


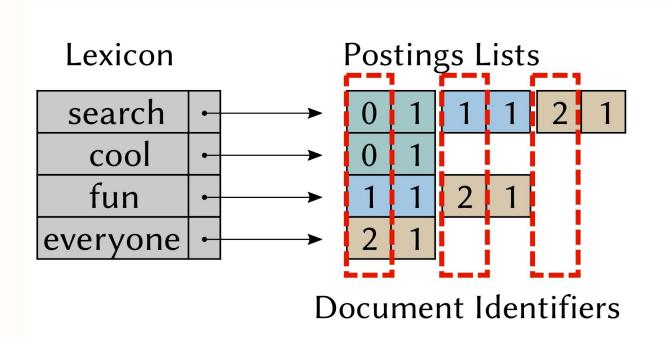


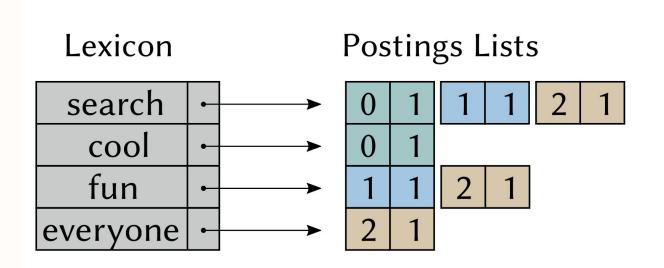


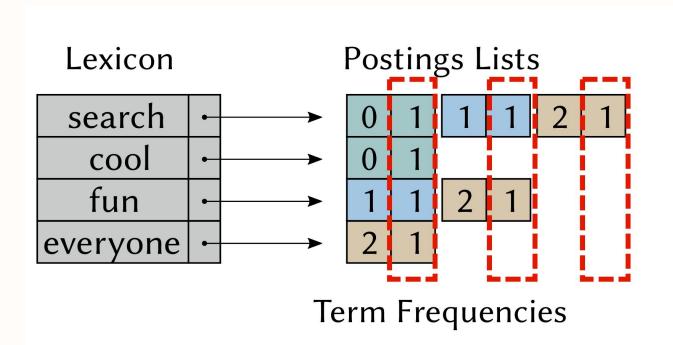


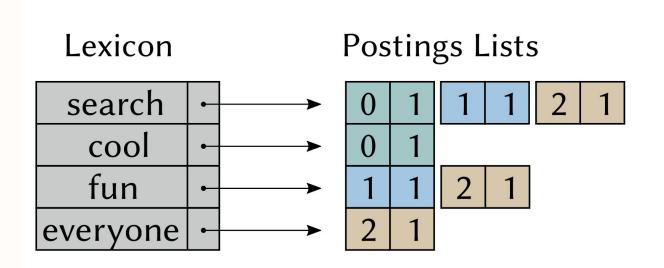


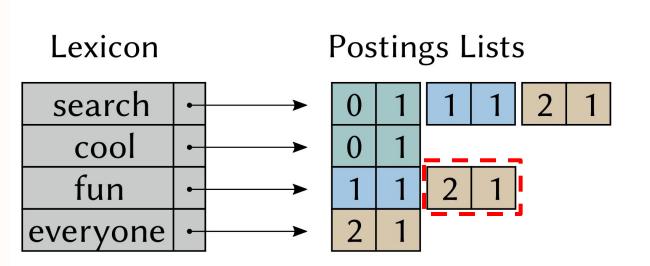




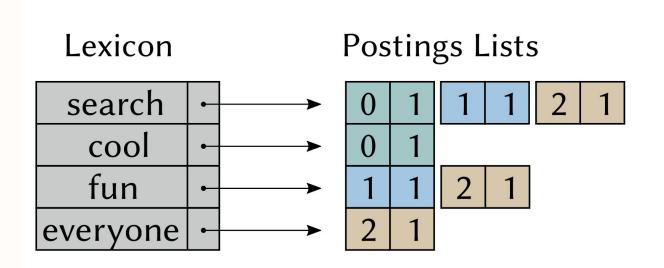


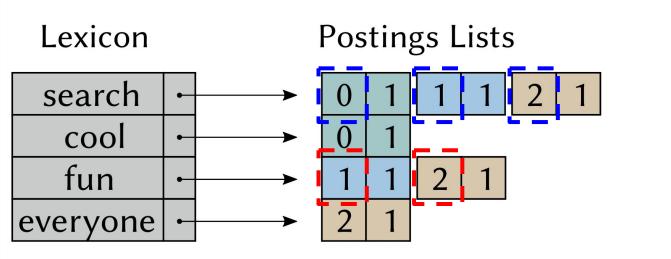






A single posting is a **document identifier** and **term frequency pair**. Document 2 contains the word "fun" once.





Within a given postings list, document identifiers are strictly increasing.

Storing Postings in Practice



Storing Postings in Practice



Storing Postings in Practice

Storing Postings in Practice

Storing Postings in Practice

Storing Postings in Practice



Storing Postings in Practice

Storing Postings in Practice



Storing Postings in Practice



Frequencies 3 4 3 1 1 5 2 3 ·

This is called **delta coding** and is a common pre-processing step used to make integer compression codecs better.

While out of scope for today, integer codecs are typically more effective for smaller integers.

Storing Postings in Practice



Storing Postings in Practice



Frequencies

Storing Postings in Practice



Frequencies



But what if I want to access the *n*th document identifier?







Storing Postings in Practice

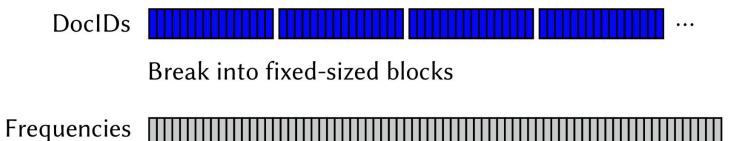
But what if I want to access the *n*th document identifier?



Frequencies

Storing Postings in Practice

But what if I want to access the *n*th document identifier?



Storing Postings in Practice

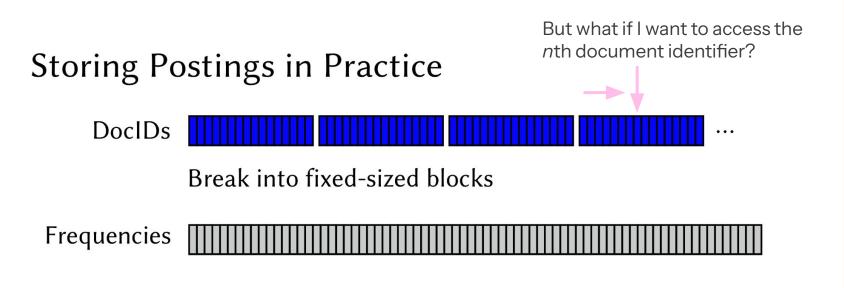
But what if I want to access the nth document identifier?



Break into fixed-sized blocks

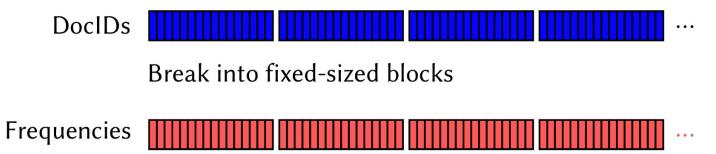


Store the *first* identifier of each *block* directly - then prefix sums are only required within each block!

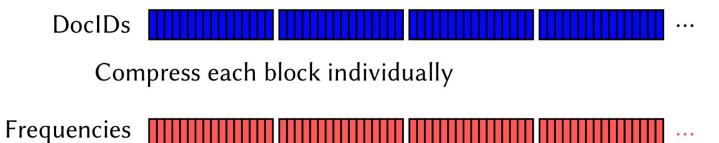


Store the *first* identifier of each *block* directly - then prefix sums are only required within each block!

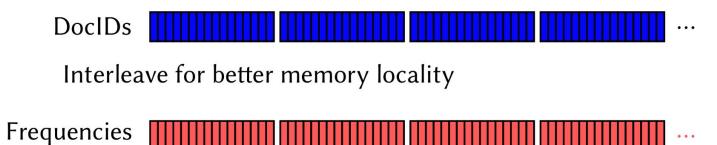
Storing Postings in Practice



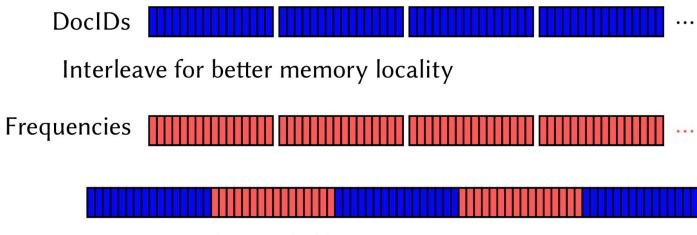
Storing Postings in Practice



Storing Postings in Practice



Storing Postings in Practice



In-memory (or on-disk) postings

Postings List Compression

Table 1. Timeline of techniques.

1949	Shannon-Fano [32, 93]
1952	Huffman [43]
1963	Arithmetic [1] ¹
1966	Golomb [40]
1971	Elias-Fano [30, 33]; Rice [87]
1972	Variable-Byte and Nibble [101]
1975	Gamma and Delta [31]
1978	Exponential Golomb [99]
1985	Fibonacci-based [6, 37]
1986	Hierarchical bit-vectors [35]
1988	Based on Front Coding [16]
1996	Interpolative [65, 66]
1998	Frame-of-Reference (For) [39]; modified Rice [2]
2003	SC-dense [11]
2004	Zeta [8, 9]

```
Simple-9, Relative-10, and Carryover-12 [3];
2005
      RBUC [60]
     PForDelta [114]; BASC [61]
      Simple-16 [112]; Tournament [100]
      ANS [27]; Varint-GB [23]; Opt-PFor [111]
      Simple8b [4]; VSE [96]; SIMD-Gamma [91]
      Varint-G8IU [97]; Parallel-PFor [5]
2013 DAC [12]; Quasi-Succinct [107]
      partitioned Elias-Fano [73]; QMX [103];
      Roaring [15, 51, 53]
      BP32, SIMD-BP128, and SIMD-FastPFor [50];
2015
      Masked-VByte [84]
2017 clustered Elias-Fano [80]
      Stream-VByte [52]; ANS-based [63, 64];
2018 Opt-VByte [83]; SIMD-Delta [104];
      general-purpose compression libraries [77];
2019 DINT [79]; Slicing [78]
```

Techniques for Inverted Index Compression - Pibiri and Venturini, ACM Computing Surveys, 2020.

Postings List Compression

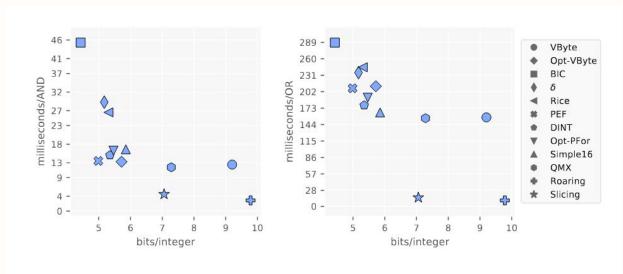


Fig. 7. Space/time trade-off curves for the ClueWeb09 dataset.

Key message: It is almost always a trade-off. Smaller codecs take longer to encode/decode, but save space.

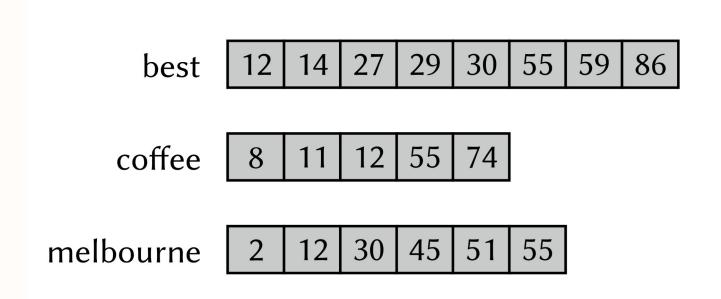
The "right" choice depends on where you wish to operate on the Pareto frontier.

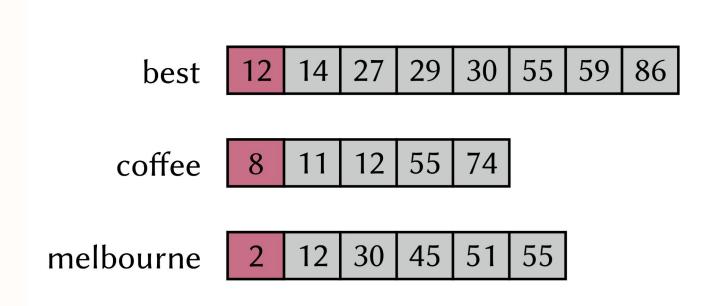
Techniques for Inverted Index Compression - Pibiri and Venturini, ACM Computing Surveys, 2020.

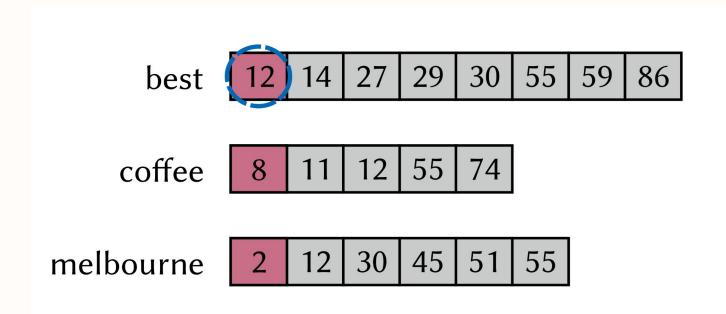
Revision: Basic Querying

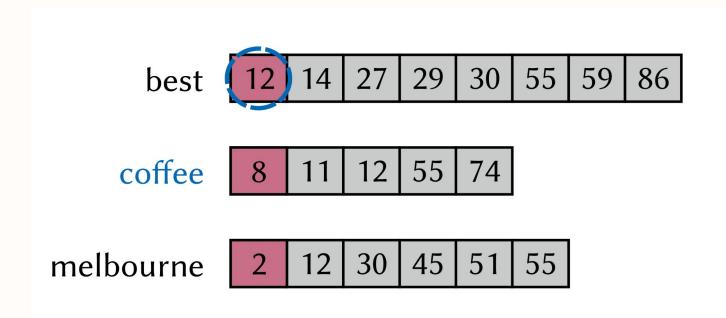
Now we have our inverted index, how can we query it?

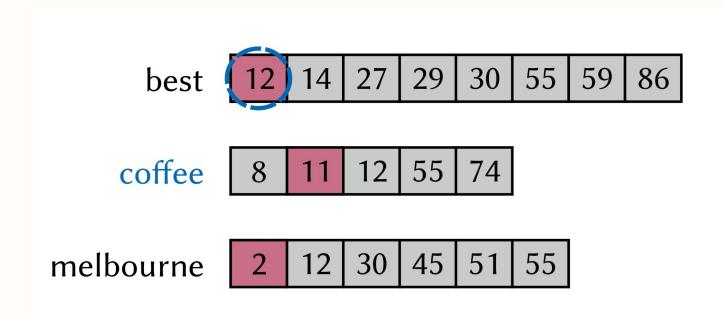
- Lots of flavours of query...
- Let's start by revising the simple Boolean conjunction
- We'll then move on to ranked disjunctions, also known as top-k retrieval.
- Important to note that *matching semantics* are separate from document *ranking*.
 - That is, we can decide to only match documents containing all query terms, but we might also decide to rank them on the way through! This would be a ranked conjunction.

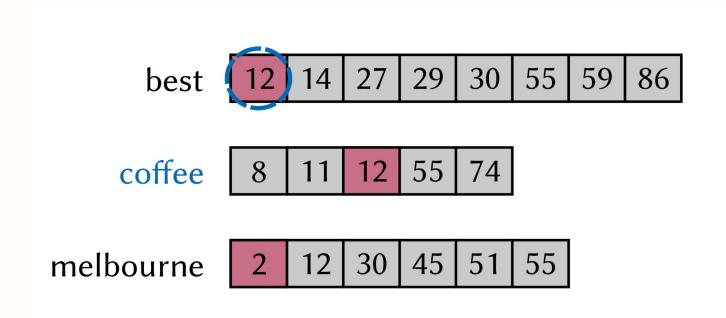


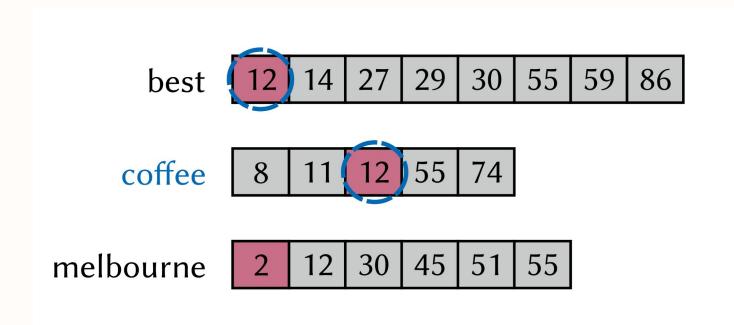


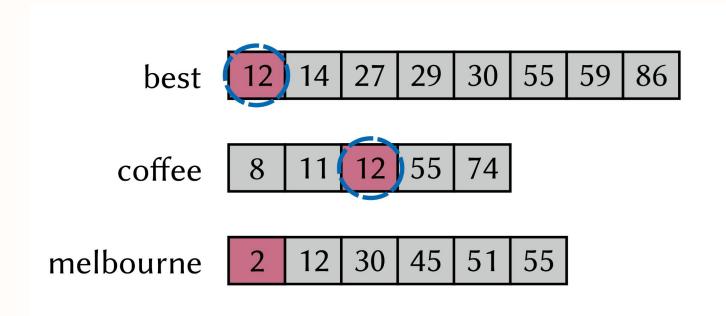


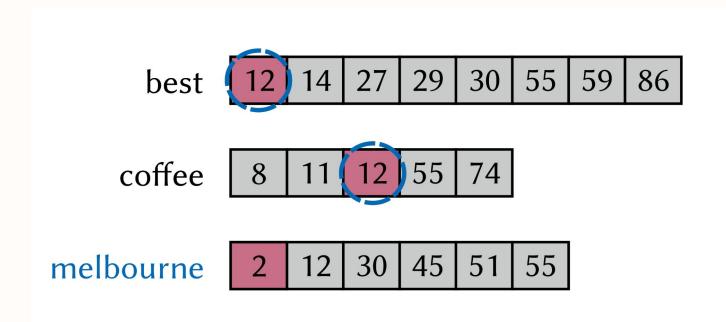


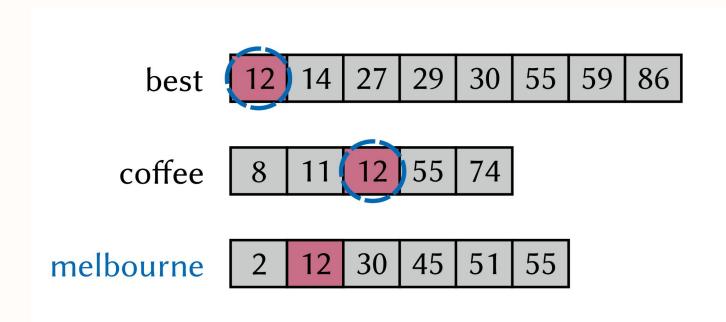


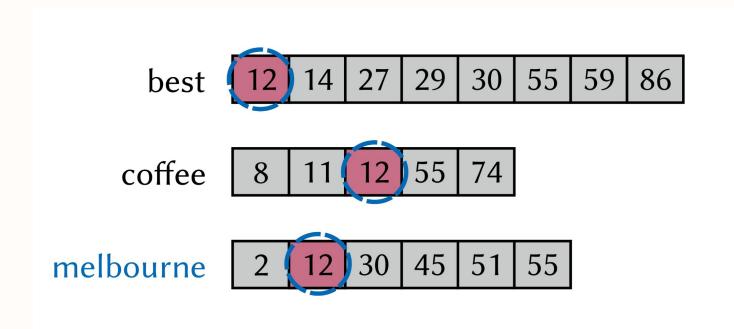


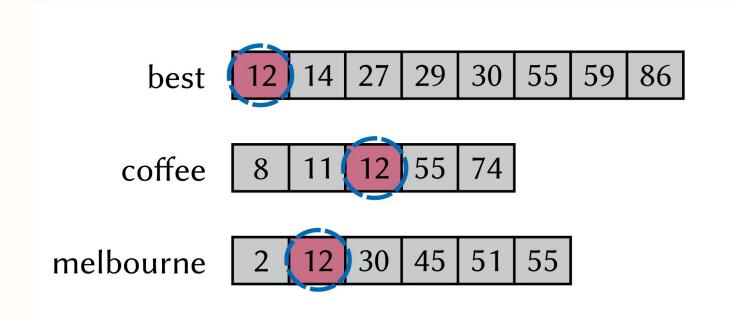


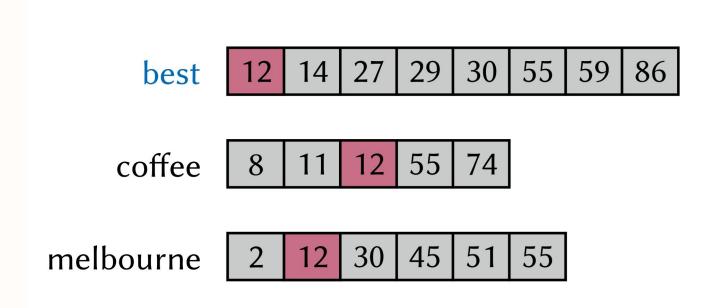


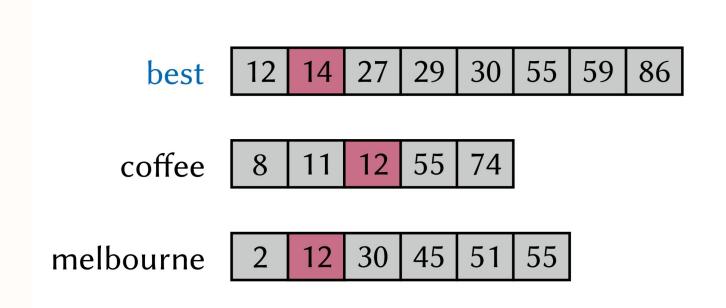


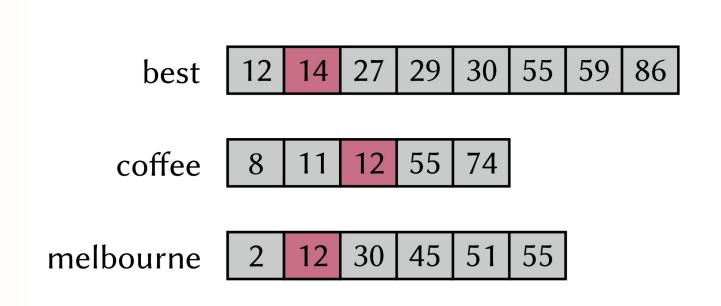


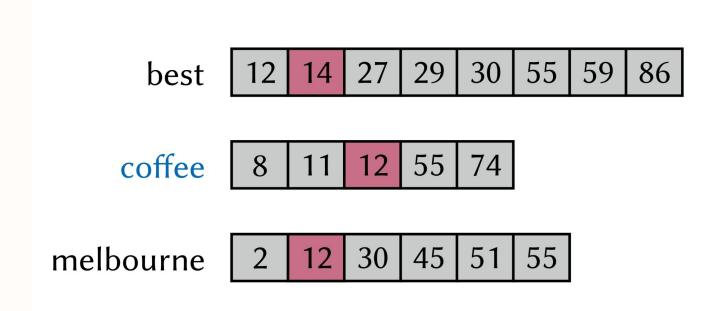


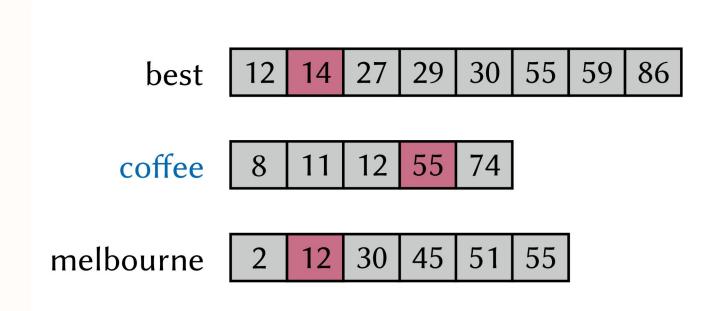


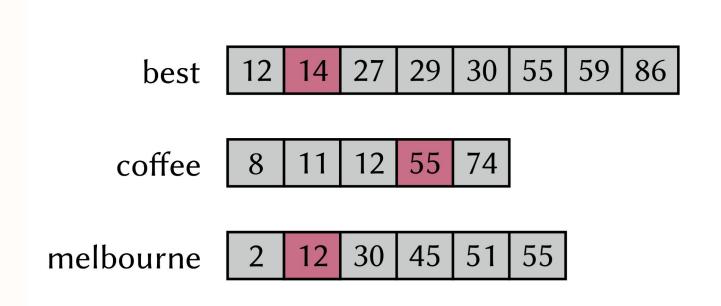


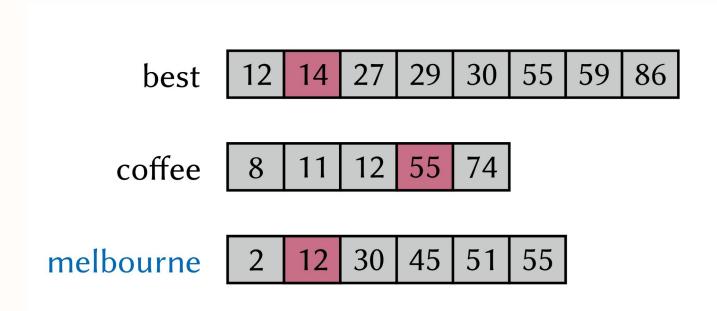


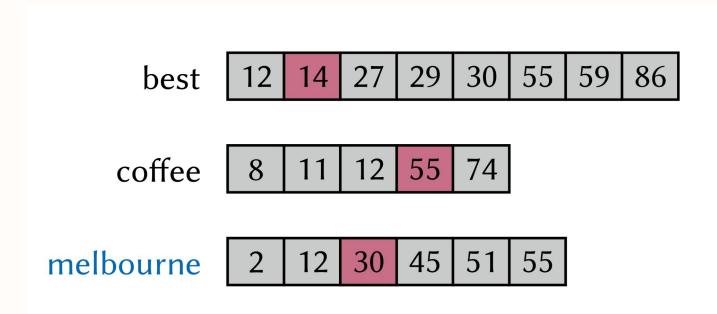


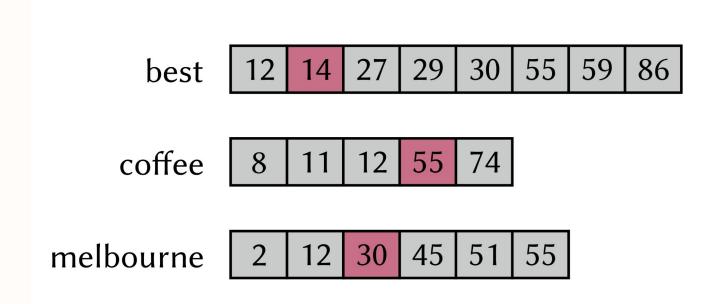


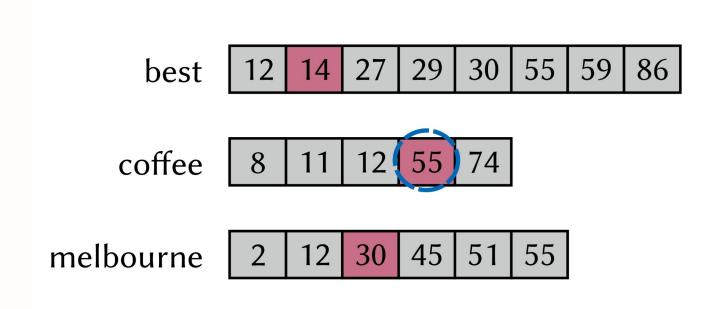


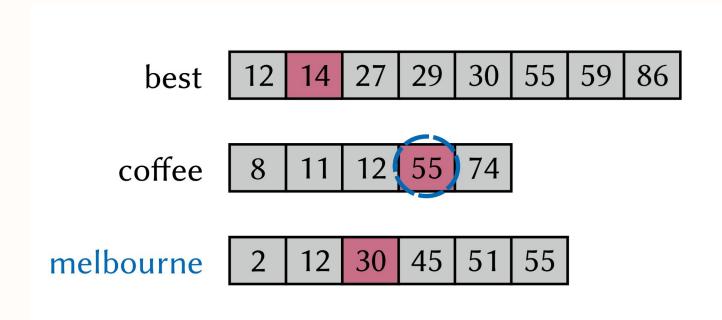


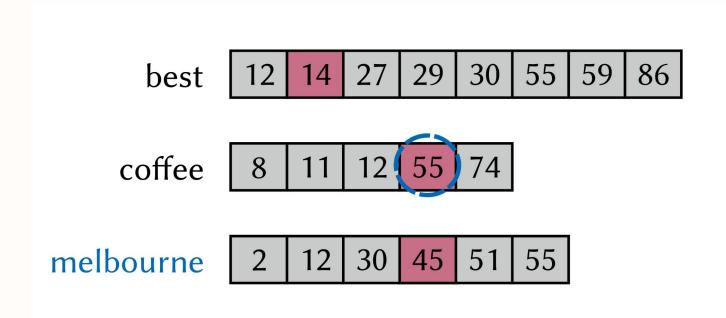


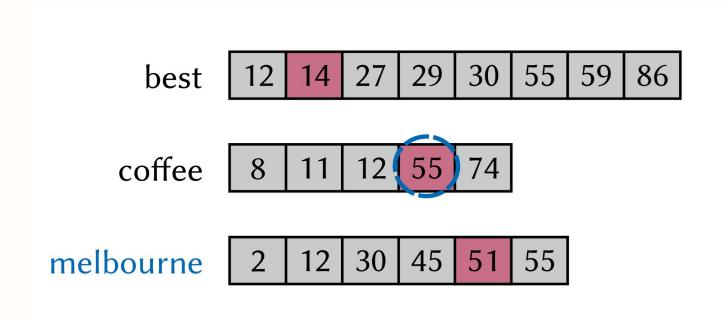


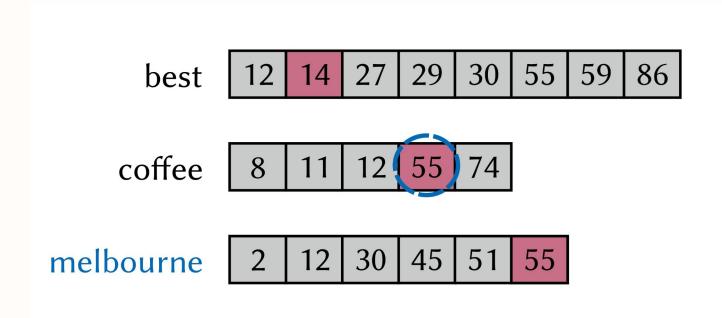


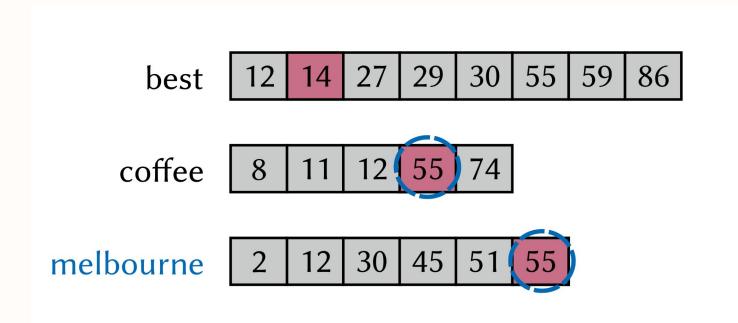


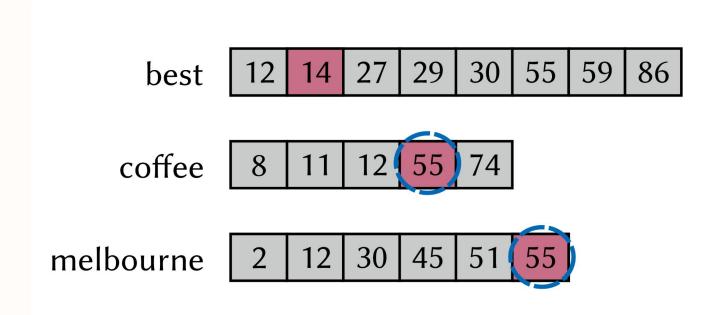


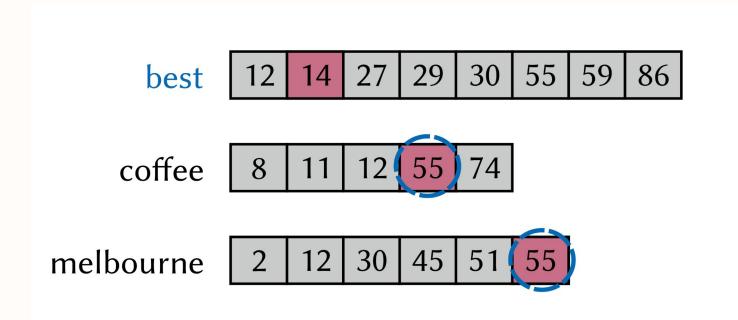


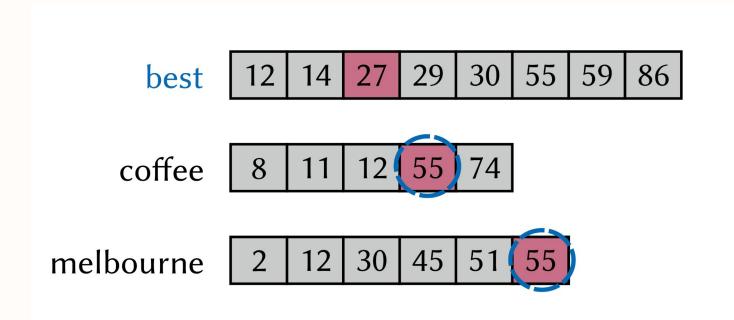


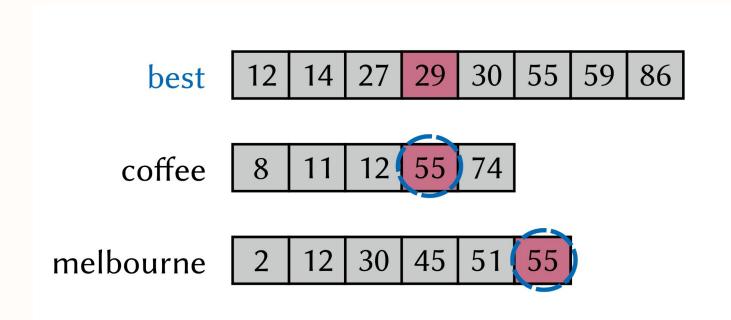


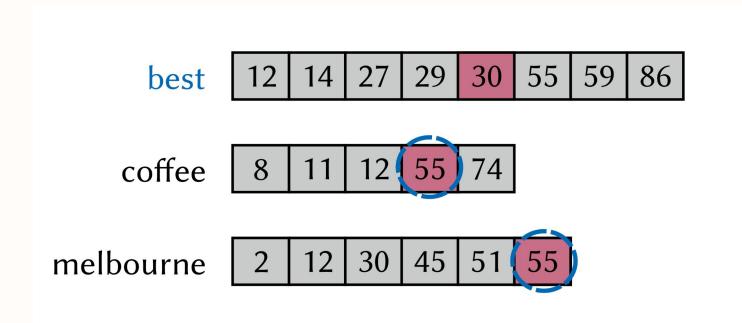


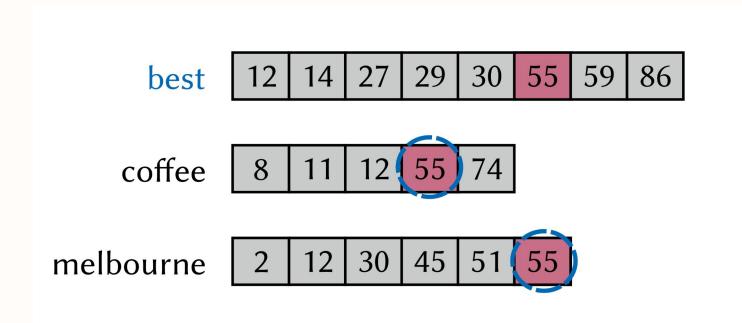


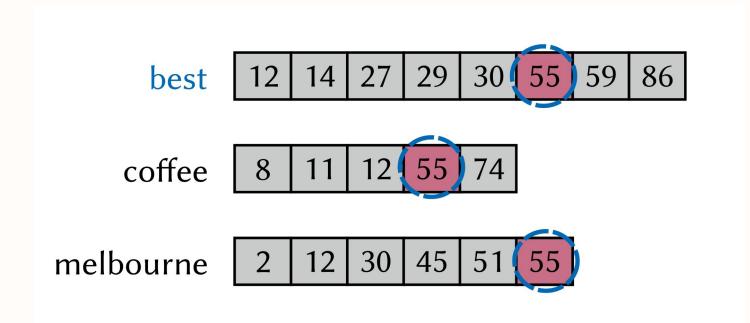




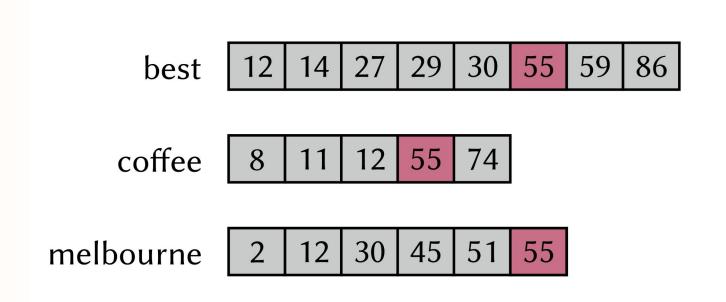






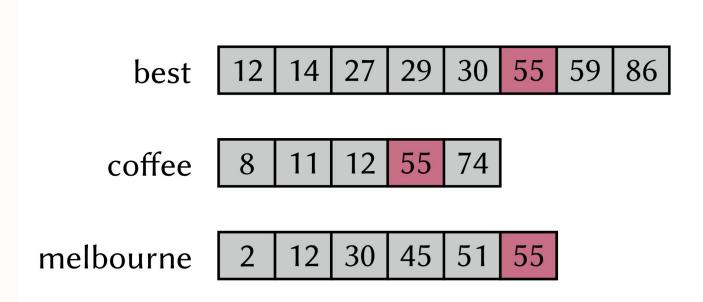


Revision: Boolean Conjunctions

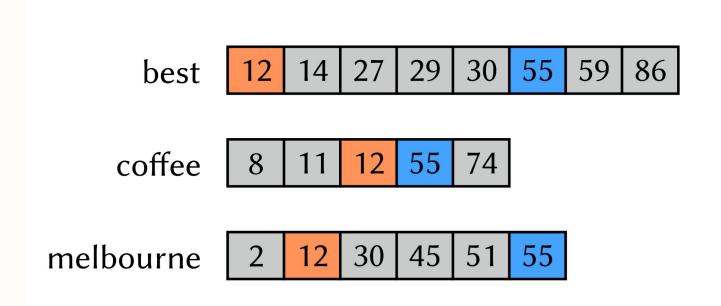


Revision: Boolean Conjunctions

And so on until we run out of postings...



Revision: Boolean Conjunctions



Consider our old, faithful friend, BM25 [ATIRE variant]:

$$\sum_{t \in q} \log \left(\frac{N}{df_t} \right) \cdot \frac{(k_1 + 1) \cdot t f_{td}}{k_1 \cdot \left(1 - b + b \cdot \left(\frac{L_d}{L_{avg}} \right) \right) + t f_{td}}$$

Consider our old, faithful friend, BM25 [ATIRE variant]:

Number of documents in the collection

Number of times t appears in "this" document

For each query term...

$$\sum\nolimits_{t \in q}^{\text{dery term...}} \log \left(\frac{N}{df_t} \right) \cdot \frac{(k_1 + 1) \cdot t f_{td}}{k_1 \cdot \left(1 - b + b \cdot \left(\frac{L_d}{L_{avg}} \right) \right) + t f_{td}}$$

Length of "this" document / Average document length

Consider our old, faithful friend, BM25 [ATIRE variant]:

$$\sum_{t \in q} \log \left(\frac{N}{df_t} \right) \cdot \frac{(k_1 + 1) \cdot t f_{td}}{k_1 \cdot \left(1 - b + b \cdot \left(\frac{L_d}{L_{avg}} \right) \right) + t f_{td}}$$

Where is this stuff stored?

Number of documents in the collection

Number of times t appears in "this" document

For each query term...

$$\sum\nolimits_{t \in q}^{\text{dery term...}} \log \left(\frac{N}{df_t}\right) \cdot \frac{(k_1 + 1) \cdot t f_{td}}{k_1 \cdot \left(1 - b + b \cdot \left(\frac{L_d}{Lavg}\right)\right) + t f_{td}}$$

Length of "this" document / Average document length

Where is this stuff stored?

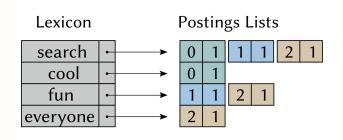
Number of documents in the collection

Number of times t appears in "this" document

 $\sum_{t \in q} \log \left(\frac{N}{df_t} \right) \cdot \frac{(k_1 + 1) \cdot t f_{td}}{k_1 \cdot \left(1 - b + b \cdot \left(\frac{L_d}{L_{ava}} \right) \right) + t f_{td}}$

Number of docs containing term t.

Length of "this" document / Average document length



+ Array with (normalized) document lengths

Consider our old, faithful friend, BM25 [ATIRE variant]:

$$\sum_{t \in q} \log \left(\frac{N}{df_t} \right) \cdot \frac{(k_1 + 1) \cdot t f_{td}}{k_1 \cdot \left(1 - b + b \cdot \left(\frac{L_d}{L_{avg}} \right) \right) + t f_{td}}$$

Note how the "query dependent" aspect is just which terms get used...

Consider our old, faithful friend, BM25 [ATIRE variant]:

$$\sum_{t \in q} \log \left(\frac{N}{df_t} \right) \cdot \frac{(k_1 + 1) \cdot t f_{td}}{k_1 \cdot \left(1 - b + b \cdot \left(\frac{L_d}{L_{avg}} \right) \right) + t f_{td}}$$

Note how the "query dependent" aspect is just which terms get used...

So, we can pre-compute the document/term impacts!

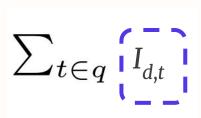
Consider our old, faithful friend, BM25 [ATIRE variant]:

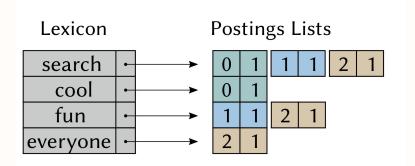
$$\sum_{t \in q} \left[I_{d,t} \right]$$

Note how the "query dependent" aspect is just which terms get used...

So, we can pre-compute the **document/term impacts!**

Consider our old, faithful friend, BM25 [ATIRE variant]:

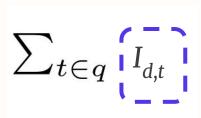


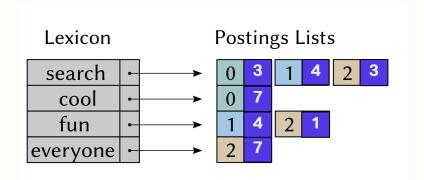


Note how the "query dependent" aspect is just which terms get used...

So, we can pre-compute the **document/term impacts**! And **store them directly in the index!**

Quantized Scoring!



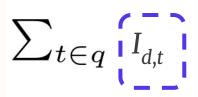


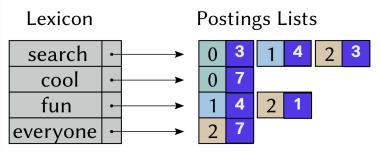
Note how the "query dependent" aspect is just which terms get used...

So, we can pre-compute the **document/term impacts!** And **store them directly in the index!** [After global normalization]

Quantized Scorers

For the remainder of this tutorial, we will assume "sum of impact" scoring



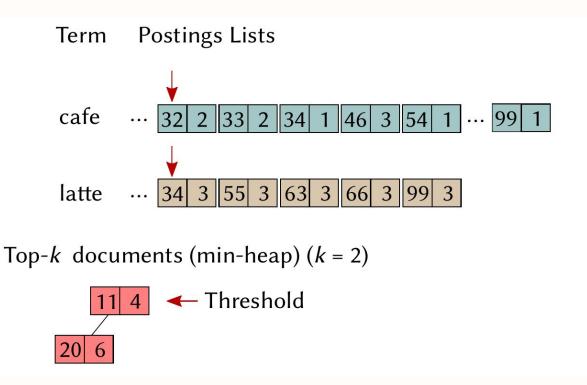


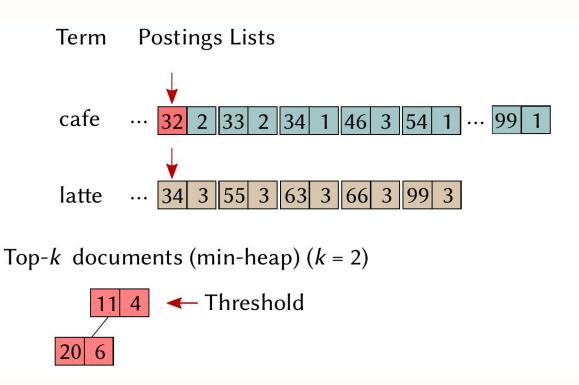
Instead of storing term frequencies, our postings lists will store quantized impacts.

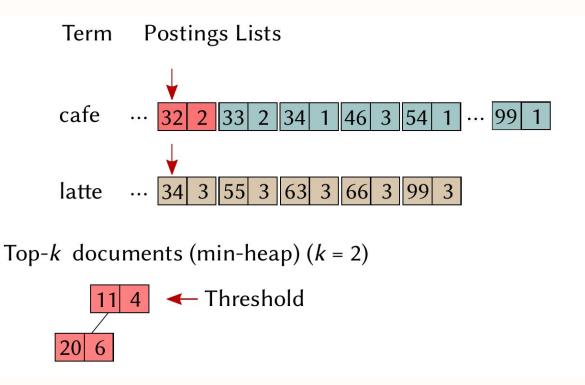
Instead of relying on conjunctive matching, let's trust our ranking function and allow *disjunctive* matching.

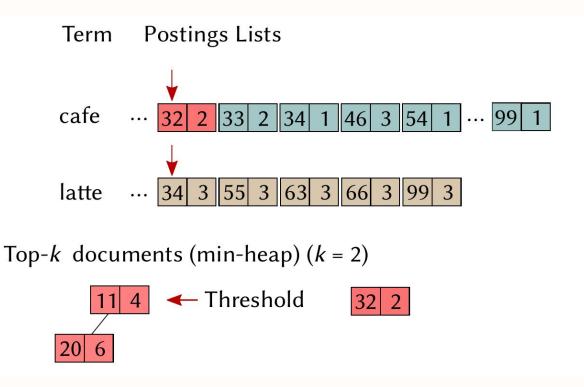
Naive algorithm: Scan across all postings lists maintaining a min-heap of the *k* best "so far" documents. For each document, compute its score, and add it to the heap if it *beats* the current top element.

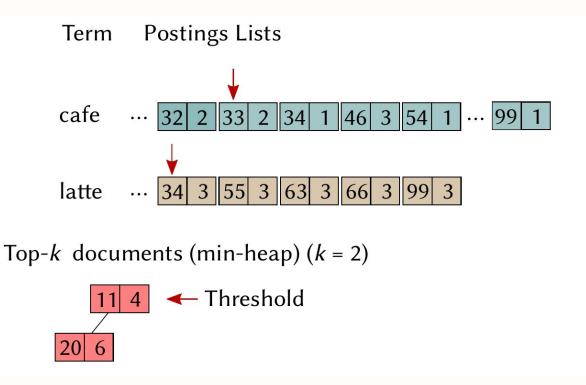
At the end of traversal, the heap contains the top-k docs.

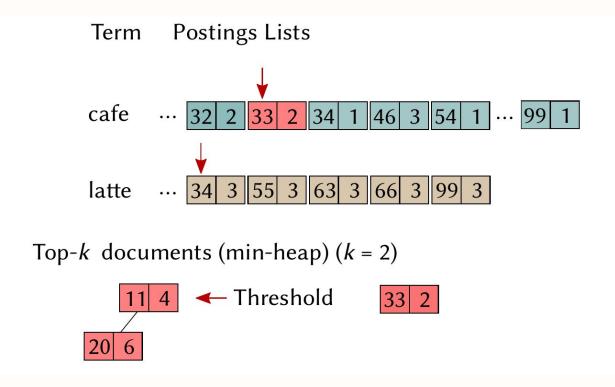


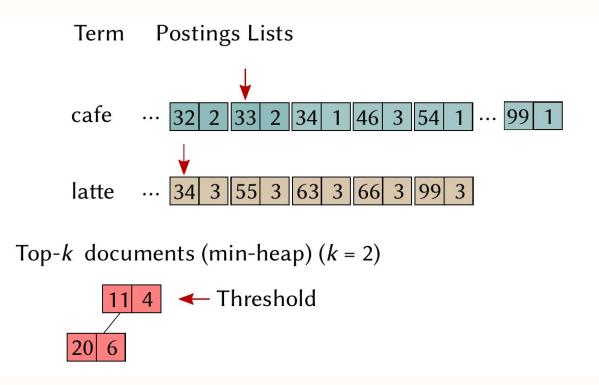


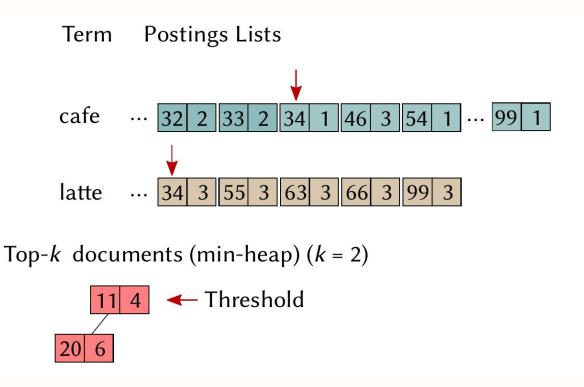


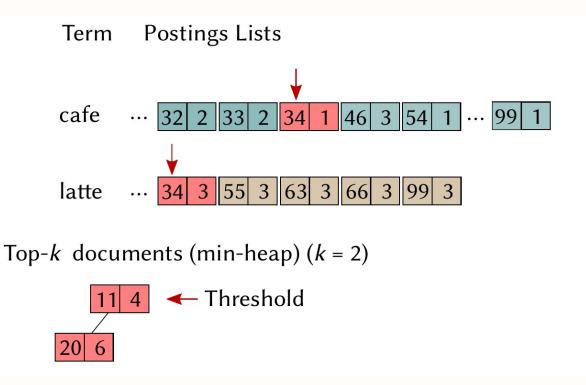


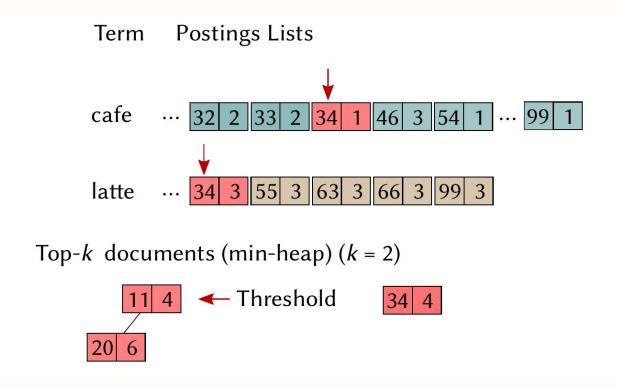


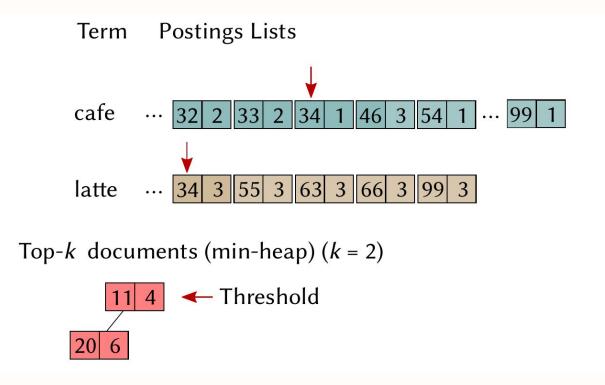


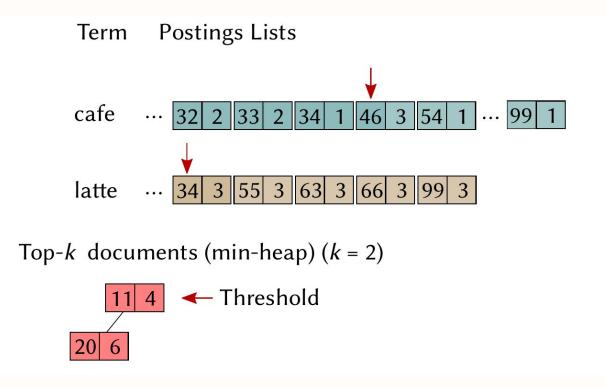


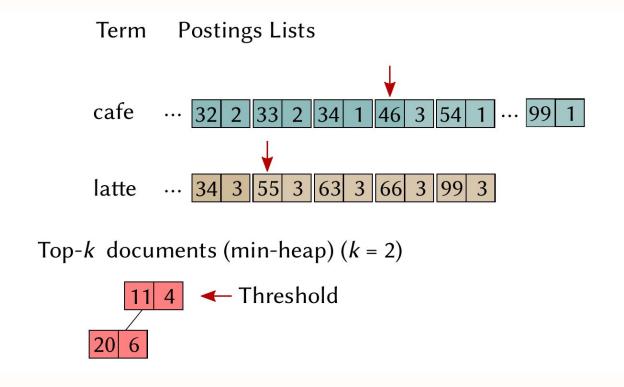












- Naive algorithm:
 - Scores every single posting!
 - Which means we decompress every block of document identifiers and impacts.
 - But guaranteed to return the rank-safe top-k documents.

Dynamic Pruning

What if we didn't need to score everything in order to get the rank-safe top-k results?

Basic Ingredients

- Our ranking function must be additive;
- We must pre-compute and store the maximum impact that each postings list contains (offline);
- Our index must support efficient random access.

Dynamic Pruning

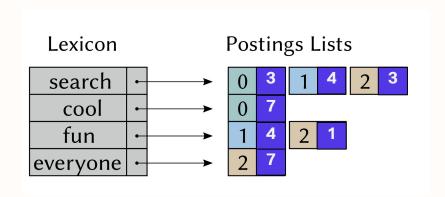
What if we didn't need to score everything in order to get the rank-safe top-k results?

Intuition

- Use the top element of the heap as a threshold;
- Estimate the score of each document by summing up the list-wise upper-bound scores;
- Only score documents with an estimated score exceeding the heap threshold - bypass otherwise.

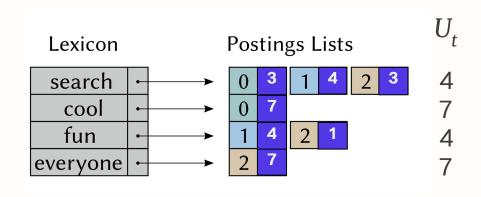
Dynamic Pruning: Index Time

• During indexing, we must pre-compute the list-wise upper-bound score. This is denoted U_t for term t.

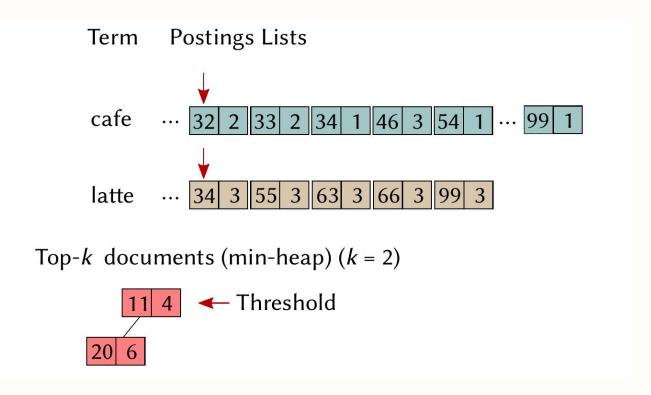


Dynamic Pruning: Index Time

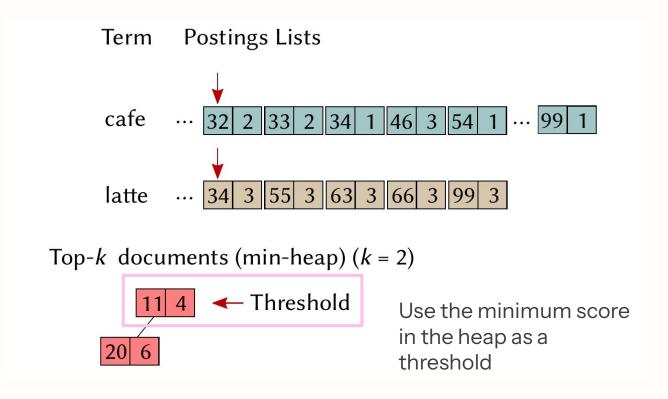
• During indexing, we must pre-compute the list-wise upper-bound score. This is denoted U_t for term t.



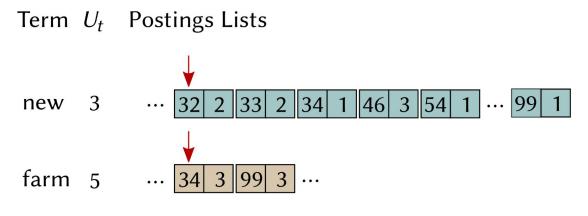
Dynamic Pruning: Query Time



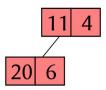
Dynamic Pruning: Query Time

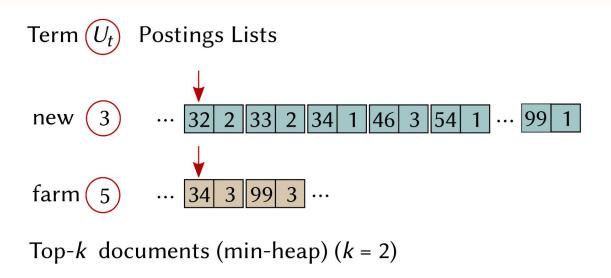


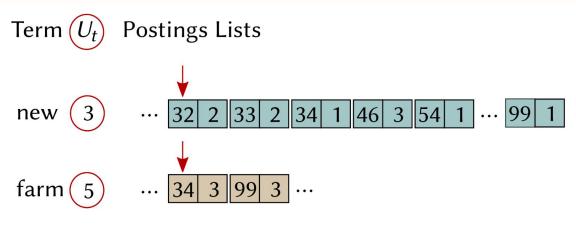
Dynamic Pruning: Walkthrough



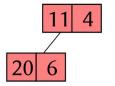
Top-k documents (min-heap) (k = 2)



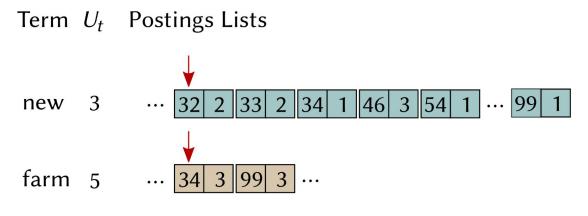




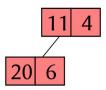
Top-k documents (min-heap) (k = 2)

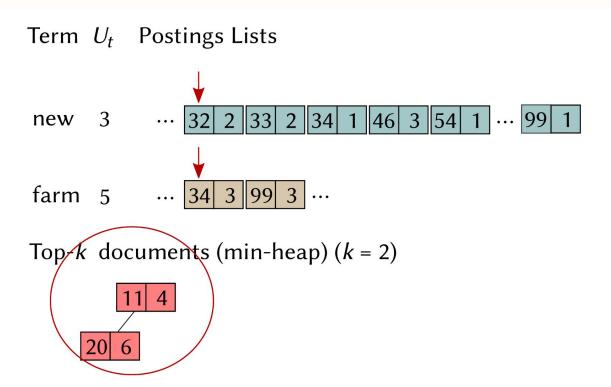


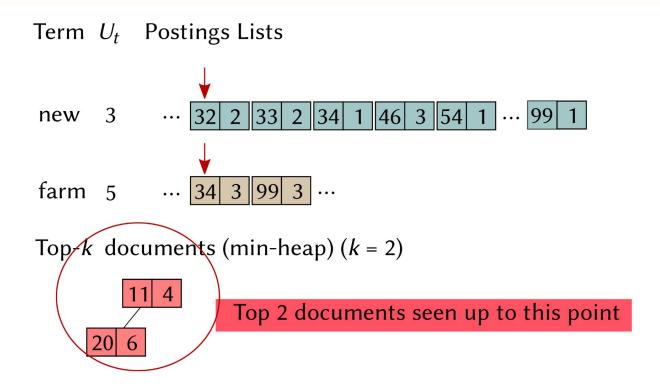
Compute and store at index time

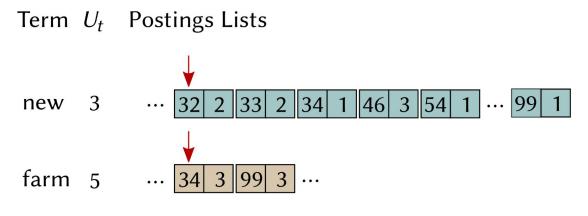


Top-k documents (min-heap) (k = 2)

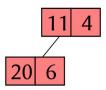


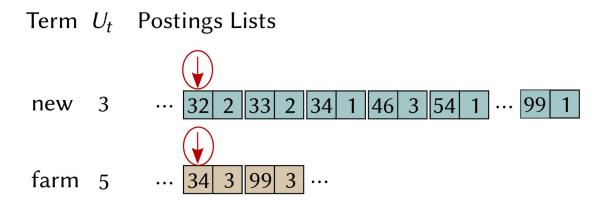




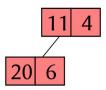


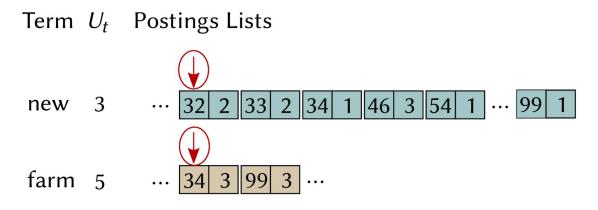
Top-k documents (min-heap) (k = 2)



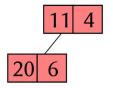


Top-k documents (min-heap) (k = 2)

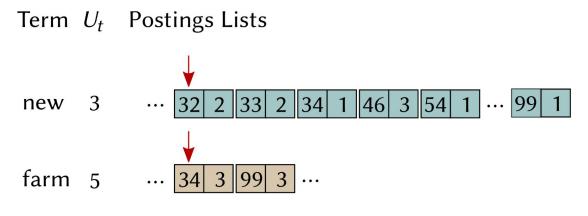




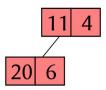
Top-k documents (min-heap) (k = 2)

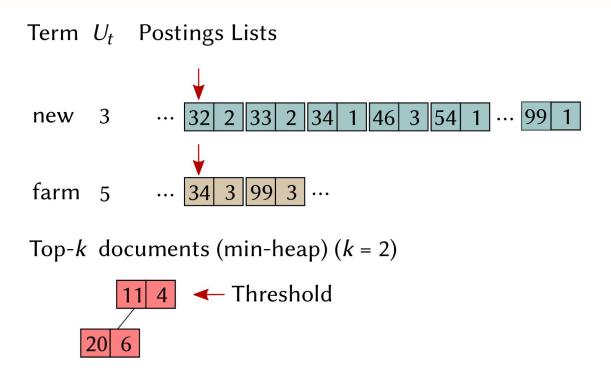


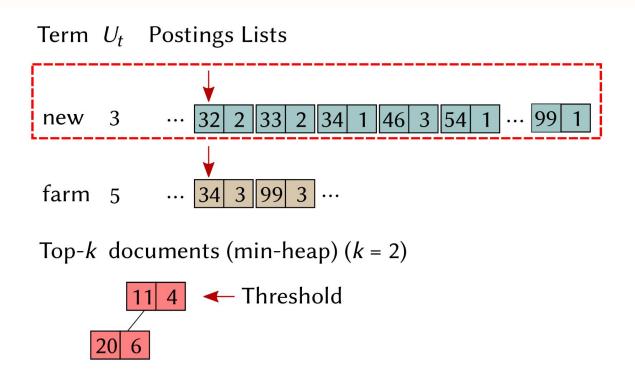
Current state of traversal

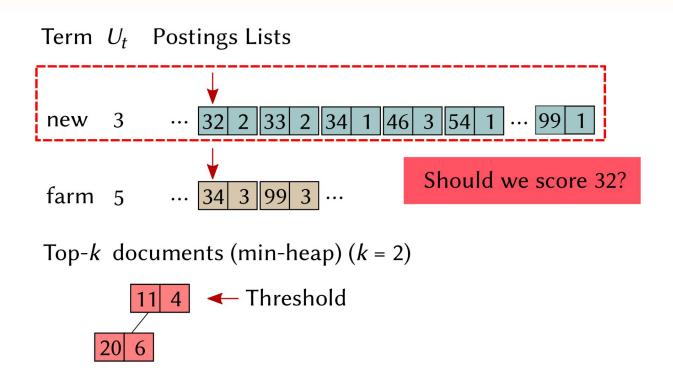


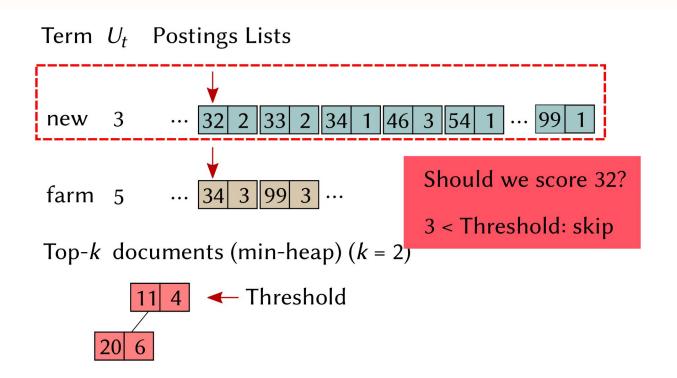
Top-k documents (min-heap) (k = 2)

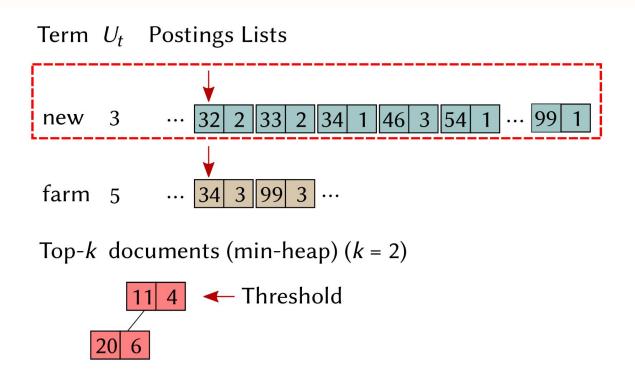


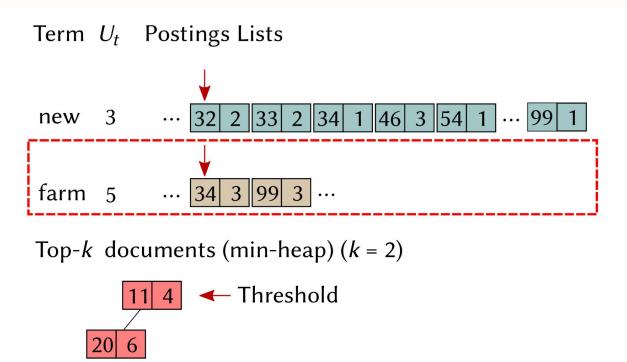


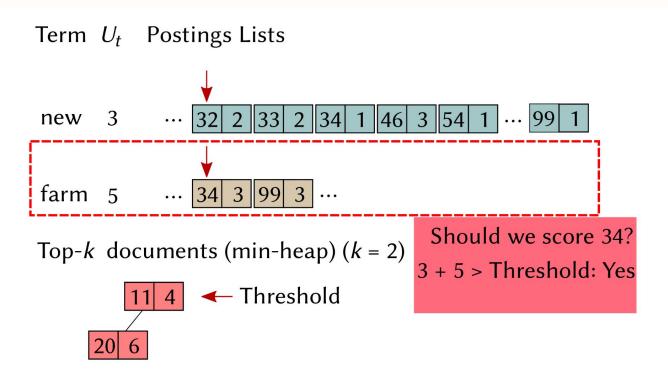


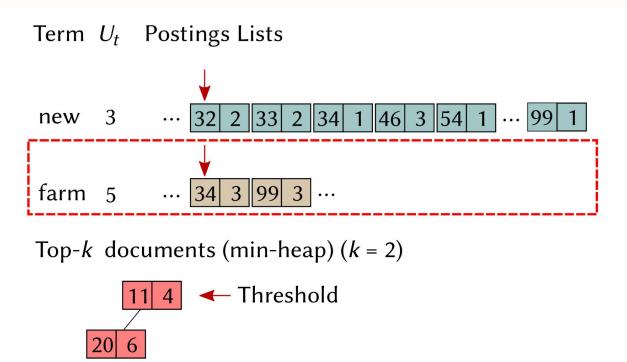


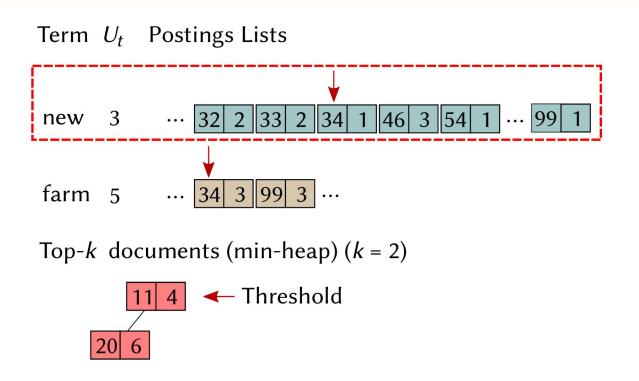


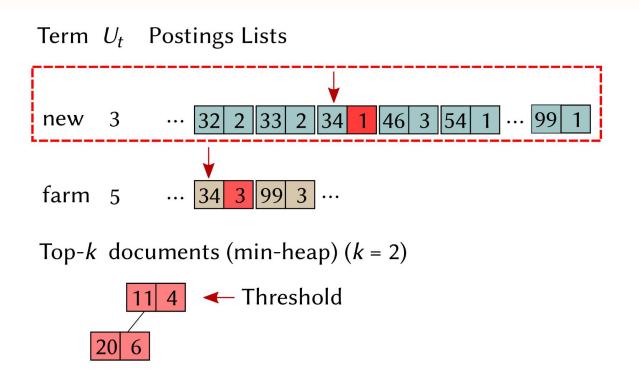


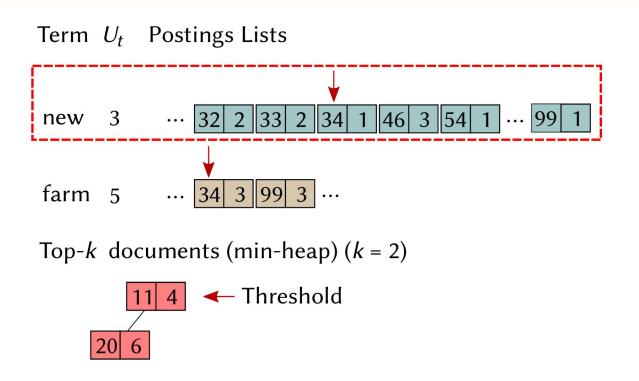


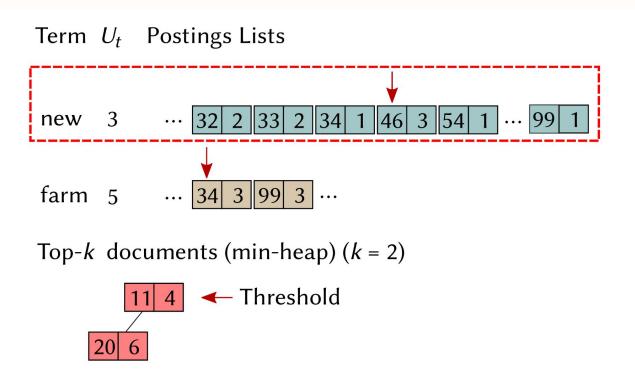






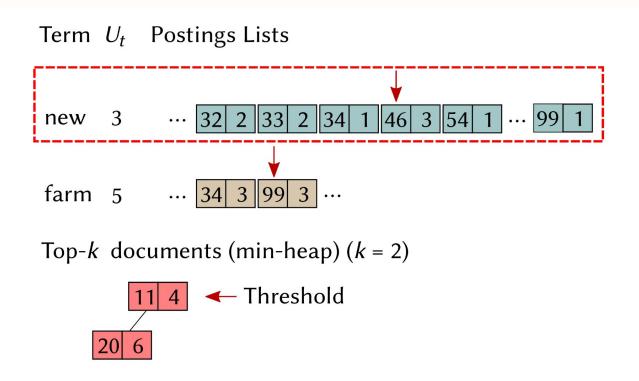


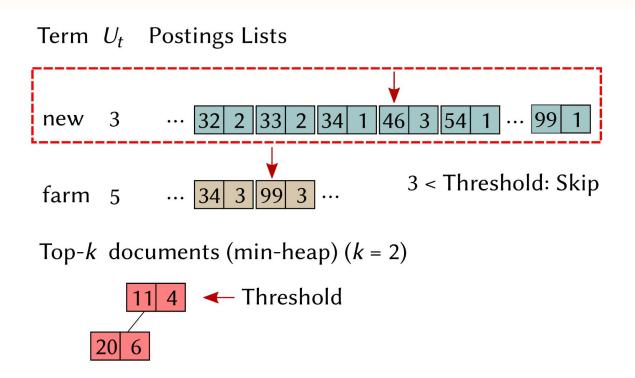


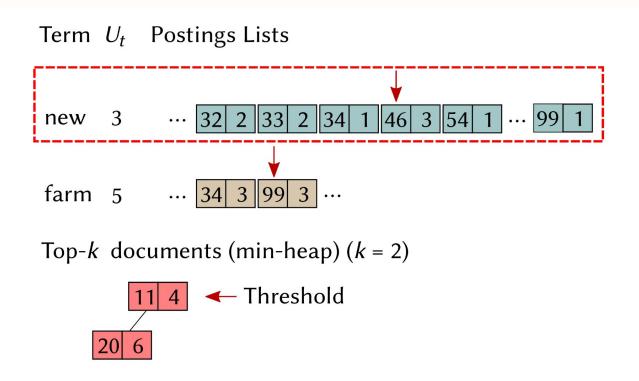


Term U_t Postings Lists ... |32| 2 ||33| 2 ||34| new farm 5 Top-k documents (min-heap) (k = 2) Threshold

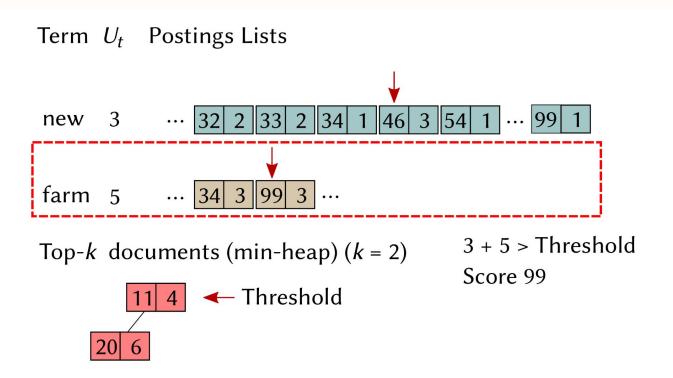
Term U_t Postings Lists ... |32| 2 ||33| 2 ||34| new farm 5 Top-k documents (min-heap) (k = 2) Threshold

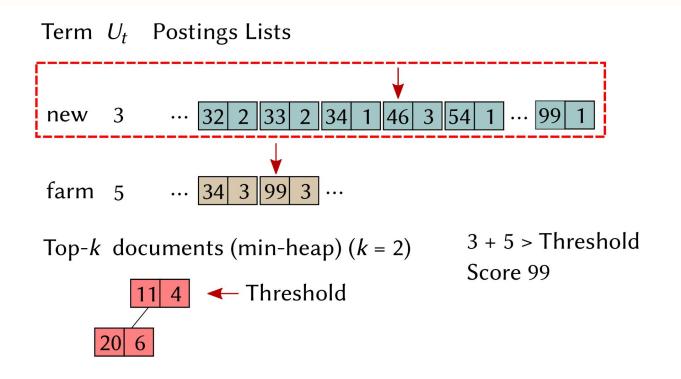


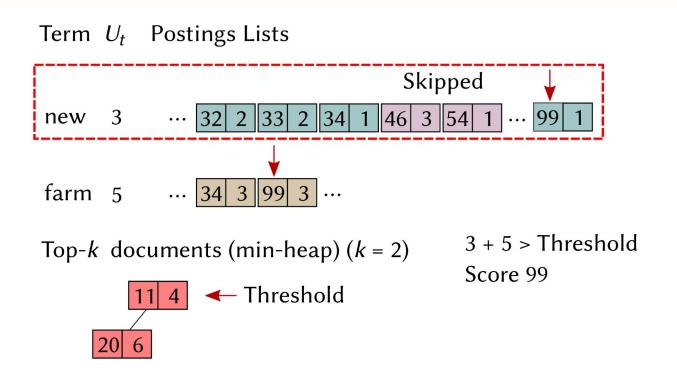


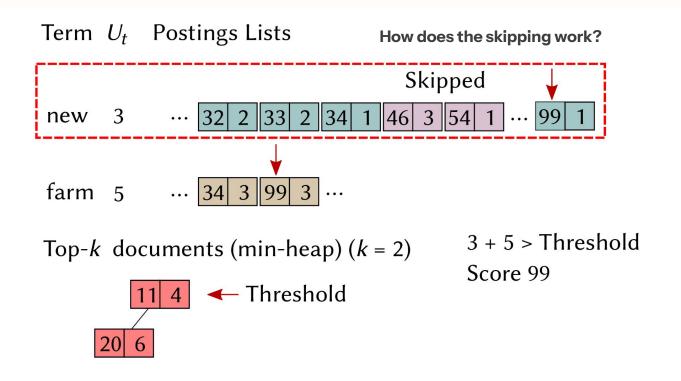


Term U_t Postings Lists ... |32| 2 ||33| 2 ||34| new farm 5 Top-k documents (min-heap) (k = 2) Threshold









Dynamic pruning algorithms depend on the **efficient** implementation of the $next_geq()$ operator.

 $next_geq(n)$ forwards the postings list cursor to document n, if it exists, or the next greater document.

In our block-based index, we retain an uncompressed document identifier such that we can skip to the candidate block efficiently; then we must seek within the block.

pisa/include/pisa/block_posting_list.hpp - Line 127

```
* Moves to the next document, counting from the current position,
* with the ID equal to or greater than `lower bound`.
 * In particular, if called with a value that is less than or equal
* to the current document ID, the position will not change.
void PISA ALWAYSINLINE next geq(uint64 t lower bound) {
   if PISA UNLIKELY (lower bound > m cur block max) {
       // binary search seems to perform worse here
       if (lower bound > block max(m blocks - 1)) {
            m cur docid = m universe;
            return:
       uint64 t block = m cur block + 1;
       while (block max(block) < lower bound) {</pre>
            ++block;
        decode docs block(block);
    while (docid() < lower bound) {</pre>
        m cur docid += m docs buf[++m pos in block] + 1;
       assert(m_pos_in_block < m_cur_block_size);
```

```
* Moves to the next document, counting from the current position,
 * with the ID equal to or greater than `lower bound`.
 * In particular, if called with a value that is less than or equal
 * to the current document ID, the position will not change.
void PISA ALWAYSINLINE next geg(uint64 t lower bound) {
   if PISA UNLIKELY (lower bound > m cur block max) {
       // pinary search seems to perform worse here
        if (lower bound > block max(m blocks - 1)) {
            m cur docid = m universe;
            return:
        uint64 t block = m cur block + 1;
        while (block max(block) < lower bound) {</pre>
            ++block;
        decode docs block(block);
    while (docid() < lower bound) {</pre>
        m cur docid += m docs buf[++m pos in block] + 1;
        assert(m_pos_in_block < m_cur_block_size);
```

If the element we are searching for is in a different block than the current one...

```
* Moves to the next document, counting from the current position,
* with the ID equal to or greater than `lower bound`.
 * In particular, if called with a value that is less than or equal
* to the current document ID, the position will not change.
void PISA ALWAYSINLINE next geq(uint64 t lower bound) {
    if PISA UNLIKELY (lower bound > m cur block max) {
        if (lower bound > block max(m blocks - 1)) {
            m cur docid = m universe;
            return:
        uint64 t block = m cur block + 1;
       while (block max(block) < lower bound) {</pre>
            ++block;
        decode_docs_block(block);
    while (docid() < lower bound) {</pre>
        m cur docid += m docs buf[++m pos in block] + 1;
       assert(m_pos_in_block < m_cur_block_size);
```

This means the element we are searching for is larger than the largest element in this postings list. So, we return.

```
* Moves to the next document, counting from the current position,
* with the ID equal to or greater than `lower bound`.
* In particular, if called with a value that is less than or equal
* to the current document ID, the position will not change.
void PISA ALWAYSINLINE next geq(uint64 t lower bound) {
   if PISA UNLIKELY (lower bound > m cur block max) {
       // binary search seems to perform worse here
       if (lower bound > block max(m blocks - 1)) {
           m cur docid = m universe;
            return:
        uint64 t block = m cur block + 1;
        while (block max(block) < lower bound) {
           ++block;
        decode docs block(block);
   while (docid() < lower bound) {</pre>
       m cur docid += m docs buf[++m pos in block] + 1;
       assert(m_pos_in_block < m_cur_block_size);
```

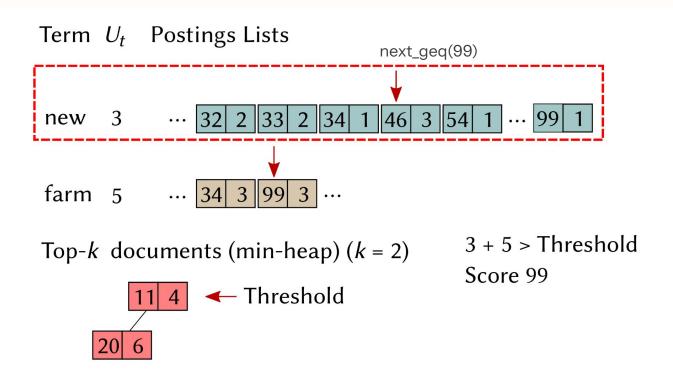
Walk across the uncompressed structure that stores the maximum identifier in each block until we exceed the target element

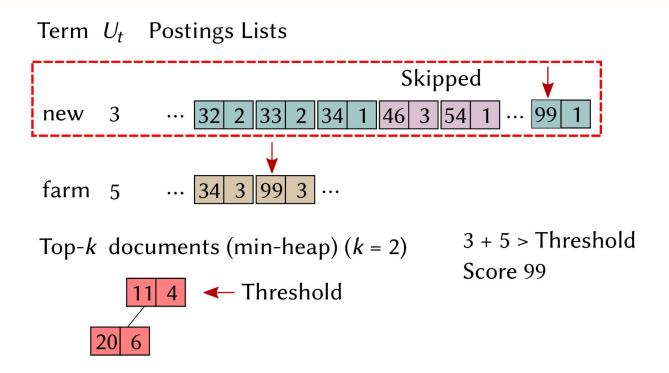
```
* Moves to the next document, counting from the current position,
* with the ID equal to or greater than `lower bound`.
* In particular, if called with a value that is less than or equal
* to the current document ID, the position will not change.
void PISA ALWAYSINLINE next geq(uint64 t lower bound) {
   if PISA UNLIKELY (lower bound > m cur block max) {
       // binary search seems to perform worse here
       if (lower bound > block max(m blocks - 1)) {
           m cur docid = m universe;
            return:
       uint64 t block = m cur block + 1;
       while (block max(block) < lower bound) {</pre>
           ++block:
        decode docs block(block);
   while (docid() < lower bound) {
       m cur docid += m docs buf[++m pos in block] + 1;
       assert(m_pos_in_block < m_cur_block_size);
```

Decode the current block into a buffer - this block must contain the target element if it exists...

```
* Moves to the next document, counting from the current position,
* with the ID equal to or greater than `lower bound`.
 * In particular, if called with a value that is less than or equal
* to the current document ID, the position will not change.
void PISA ALWAYSINLINE next geq(uint64 t lower bound) {
   if PISA UNLIKELY (lower bound > m cur block max) {
       // binary search seems to perform worse here
       if (lower bound > block max(m blocks - 1)) {
            m cur docid = m universe;
            return:
       uint64 t block = m cur block + 1;
       while (block max(block) < lower bound) {</pre>
            ++block:
        decode docs block(block);
   while (docid() < lower bound) {</pre>
        m cur docid += m docs buf[++m pos in block] + 1;
       assert(m pos in block < m cur block size);
```

Now search within the block to find the element (or the next greater one)





Skipping - Saving Work

We do not need to decompress any blocks we skip!

We do not need to score any documents within any of those blocks either!

But we do pay **overhead** deciding whether to skip or not.

We also need to be **very careful** with which cursor moves ahead first. More on this shortly...

Efficient Query Evaluation using a Two-Level Retrieval Process

Andrei Z. Broder[§], David Carmel*, Michael Herscovici*, Aya Soffer*, Jason Zien[†]

(§) IBM Watson Research Center, 19 Skyline Drive, Hawthorne, NY 10532

(*) IBM Research Lab in Haifa, MATAM, Haifa 31905, ISRAEL

(†) IBM Almaden Research Center, 650 Harry Road, San Jose, CA 95120

WAND - "Weak" or "Weighted" AND CIKM 2003!

QUERY EVALUATION: STRATEGIES AND OPTIMIZATIONS

HOWARD TURTLE and JAMES FLOOD

West Publishing Co., 610 Opperman Drive, Eagan, MN 55123, U.S.A.

(Received January 1995; accepted in final form March 1995)

Abstract—This paper discusses the two major query evaluation strategies used in large text retrieval systems and analyzes the performance of these strategies. We then discuss several optimization techniques that can be used to reduce evaluation costs and present simulation results to compare the performance of these optimization techniques when evaluating natural language queries with a collection of full text legal materials.

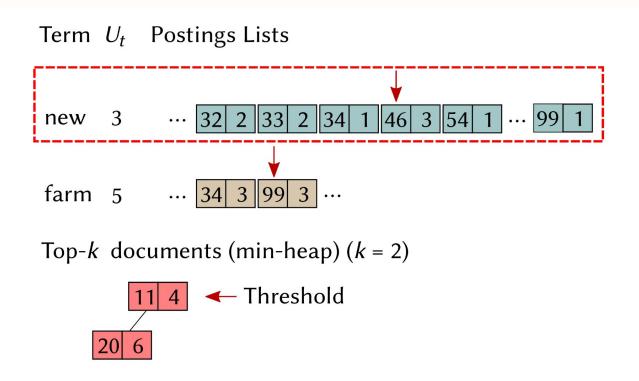
MaxScore IPM 1995!

WAND and MaxScore are the two main families of document-at-a-time dynamic pruning algorithms.

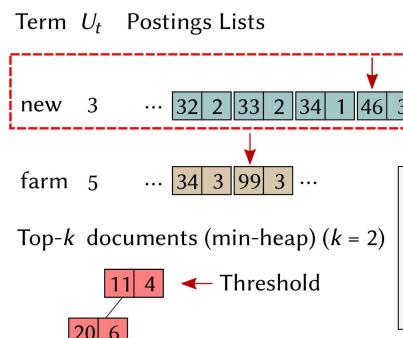
Our worked example was inspired by the WAND algorithm.

With WAND, we must ensure the cursors are always processed in ascending order of their current identifier.

Dynamic Pruning: Walkthrough



Dynamic Pruning: Walkthrough

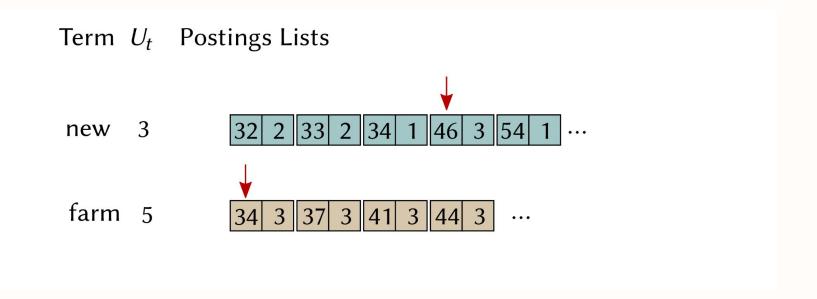


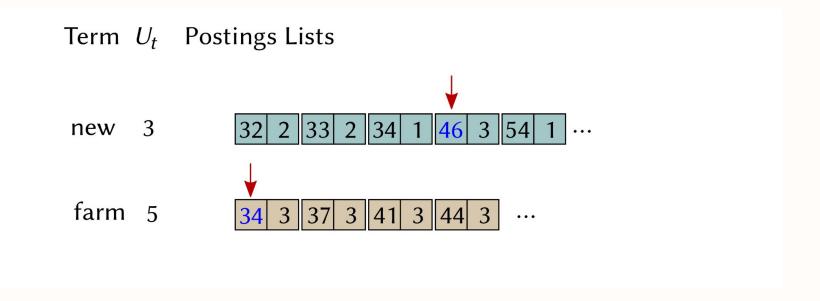
Invariant: The current identifier under each cursor must monotonically increase as we move down through the cursors.

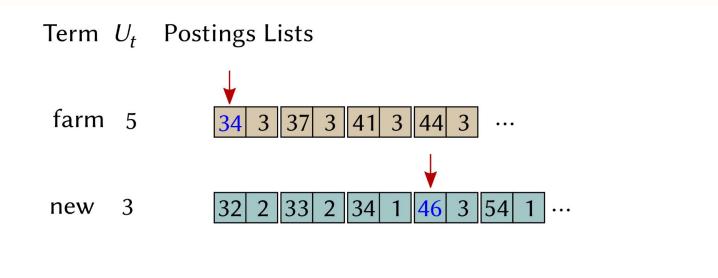
That is, "the element being pointed to in list at the top must be less than or equal to the element being pointed to in the list at the bottom"

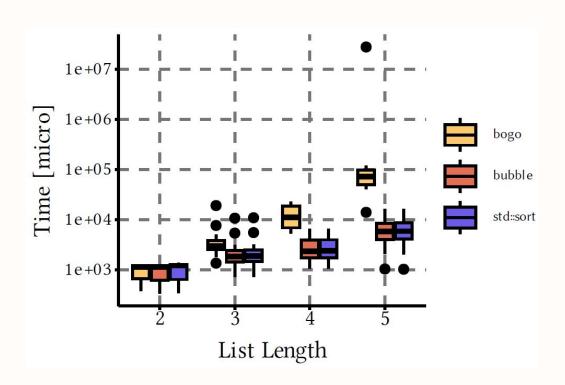
With WAND, we must ensure the cursors are always processed in ascending order of their current identifier.

This means we may need to **sort** the cursors during processing!









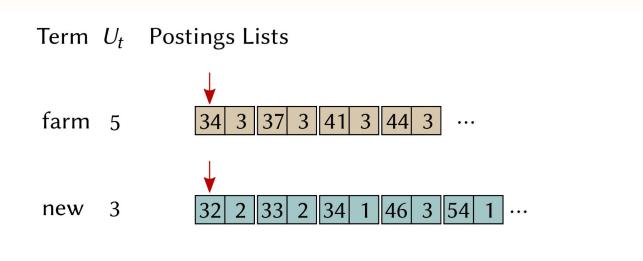
Just for fun... What happens if we change the sorting algorithm?

Instead of sorting the cursors before each iteration, MaxScore sorts the lists *once* before processing begins.

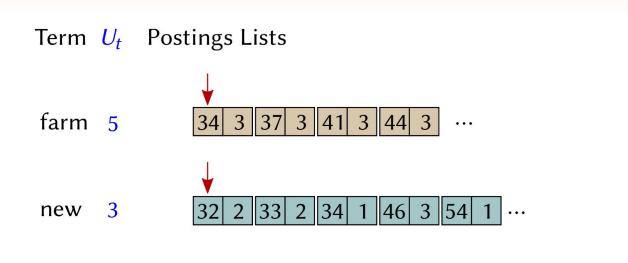
Typically, this sort is ascending on the list upper bounds.

Can also sort on the posting list lengths - either works.

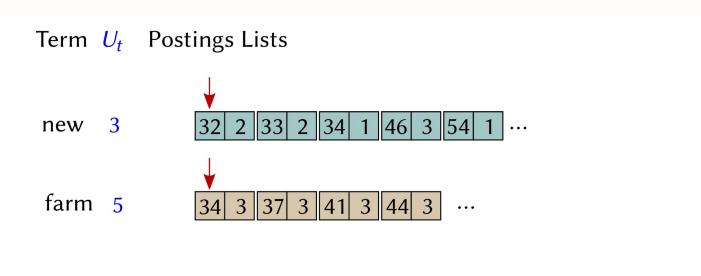
Instead of sorting the cursors before each iteration, MaxScore sorts the lists **once** before processing begins.



Instead of sorting the cursors before each iteration, MaxScore sorts the lists **once** before processing begins.

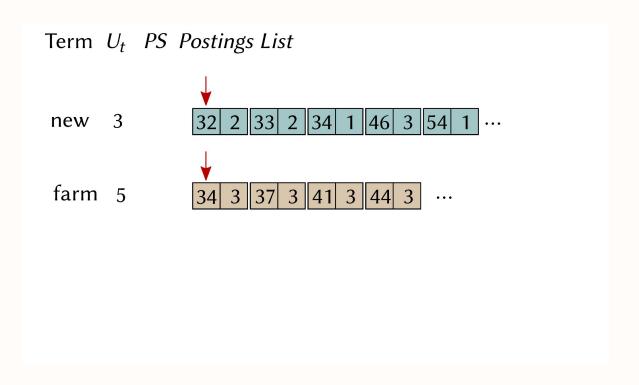


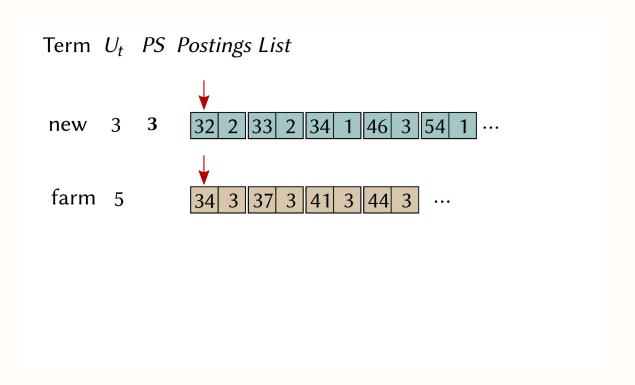
Instead of sorting the cursors before each iteration, MaxScore sorts the lists **once** before processing begins.

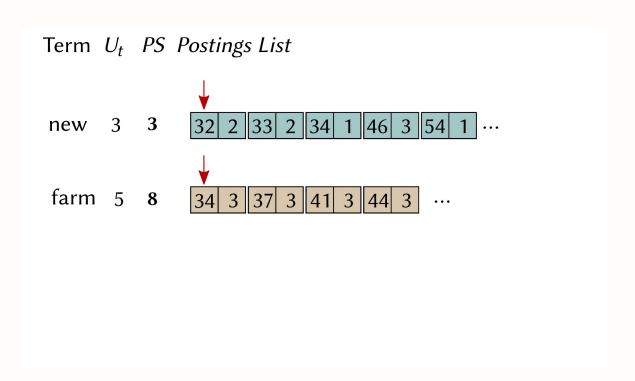


Instead of sorting the cursors before each iteration, MaxScore sorts the lists **once** before processing begins.

We then compute the *prefix sum* of the upper-bound scores.



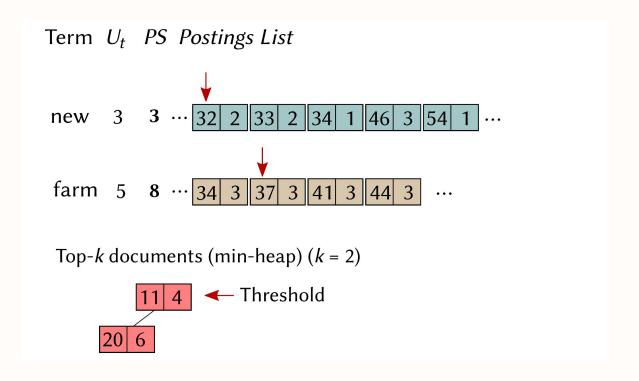


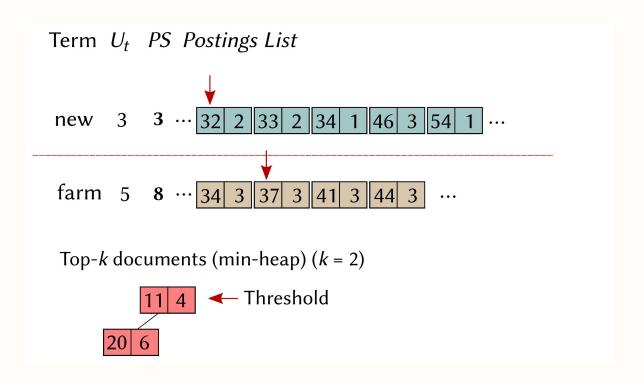


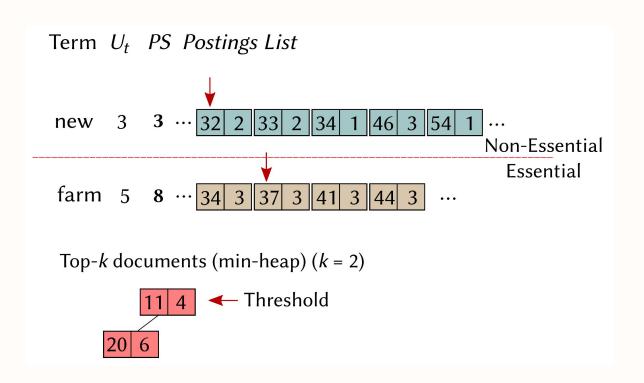
During processing, we split the cursors into two logical sets:

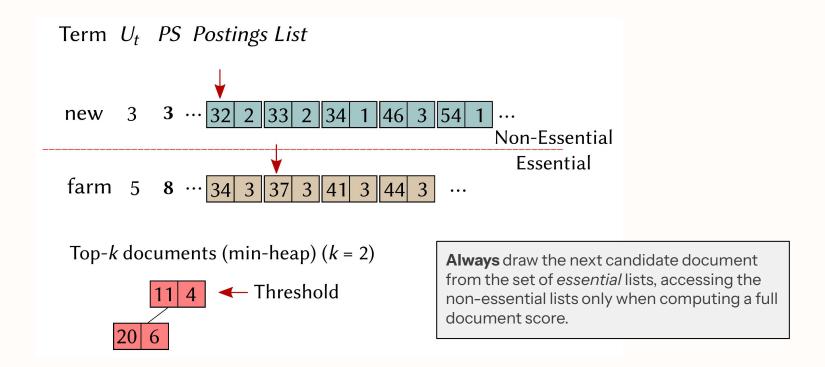
- **Essential Lists:** A document must have at least one essential term to be considered for scoring.
- Non-Essential Lists: Are ignored until a final document score needs to be computed.

We only iterate over the *essential* lists, allowing us to skip documents that have no hope of entering the results list.









Algorithm 3.3: WAND processing.

```
Input : An array \mathcal{P} of n postings cursors, the largest document identifier, MaxDocID,
            and the number of desired results, k.
  Output : The top-k documents.
 1 Heap ← {}
2 θ ← 0
3 SortByDocID(P)
 4 while true do
      UpperBound ← 0
      PivotList ← 0
      PivotID \leftarrow MaxDocID
      while PivotList < n do
                                     // Find the pivot list and pivot document.
          UpperBound \leftarrow UpperBound +\mathcal{P}[PivotList].upperbound()
          if UpperBound > \theta then
10
              PivotID \leftarrow P[PivotList].docid()
11
              while PivotList+1 < n and P[PivotList+1].docid() = PivotID do
12
                 PivotList \leftarrow PivotList +1
13
14
              end
              break
15
          end
16
          PivotList ← PivotList +1
17
18
      if UpperBound \leq \theta or PivotID = MaxDocID then
                                                                    // No pivot. Exit.
         break
20
      end
21
      if \mathcal{P}[0].docid() = PivotID then
                                             // Can evaluate the pivot document.
22
          Score ← 0
23
          for List ← 0 to PivotList do
24
              Score \leftarrow Score +\mathcal{P}[List].score()
25
              \mathcal{P}[\text{List}].next()
26
          end
27
          Heap.push((PivotID, Score))
28
          \theta \leftarrow \mathsf{Heap}.min()
29
          SortByDocID(P)
30
31
      end
      else
                                          // Need to align lists with the pivot.
32
          while \mathcal{P}[\mathsf{PivotList}] = \mathsf{PivotID} \, \mathbf{do}
33
              PivotList ← PivotList -1
34
          end
35
          P[PivotList].nextGEQ(PivotID)
36
          BubbleDown(P, PivotList)
37
      end
38
39 end
40 return Heap
```

Algorithm 3.3: WAND processing.

```
Input: An array \mathcal{P} of n postings cursors, the largest document identifier, MaxDocID,
            and the number of desired results, k.
  Output: The top-k documents.
 1 Heap ← {}
2 θ ← 0
 3 SortByDocID(P)
 4 while true do
      UpperBound ← 0
      PivotList ← 0
      PivotID - MayDocID
                                    // Find the pivot list and pivot document.
       while PivotList < n \text{ do}
          UpperBound \leftarrow UpperBound +\mathcal{P}[PivotList].upperbound()
          if UpperBound > \theta then
              PivotID \leftarrow P[PivotList].docid()
11
             while PivotList+1 < n and P[PivotList+1].docid() = PivotID do
12
                  PivotList ← PivotList +1
14
              end
15
              break
          end
16
          PivotList ← PivotList +1
18
      if UpperBound \leq \theta or PivotID = MaxDocID then
                                                                    // No pivot. Exit.
         break
      end
21
      if \mathcal{P}[0].docid() = PivotID then
22
                                             // Can evaluate the pivot document.
23
          Score \leftarrow 0
          for List ← 0 to PivotList do
24
              Score \leftarrow Score +\mathcal{P}[List].score()
25
             \mathcal{P}[\text{List}].next()
26
          end
27
          Heap.push((PivotID, Score))
28
          \theta \leftarrow \text{Heap.}min()
29
          SortByDocID(P)
30
31
      end
                                          // Need to align lists with the pivot.
32
      else
          while \mathcal{P}[PivotList] = PivotID do
33
              PivotList \leftarrow PivotList -1
34
          end
35
          P[PivotList].nextGEQ(PivotID)
36
          BubbleDown(P, PivotList)
      end
38
39 end
40 return Heap
```

Find the first document that *might* enter the top-*k* based on the upper-bound estimations

This is called the "pivot" document.

We also track which lists are "in play".

Algorithm 3.3: WAND processing.

```
Input : An array \mathcal{P} of n postings cursors, the largest document identifier, MaxDocID,
            and the number of desired results, k.
  Output : The top-k documents.
 1 Heap ← {}
2 θ ← 0
3 SortByDocID(P)
 4 while true do
      UpperBound ← 0
      PivotList ← 0
      PivotID ← MaxDocID
      while PivotList < n do
                                    // Find the pivot list and pivot document.
          UpperBound \leftarrow UpperBound +\mathcal{P}[PivotList].upperbound()
          if UpperBound > \theta then
10
             PivotID \leftarrow P[PivotList].docid()
11
             while PivotList+1 < n and \mathcal{P}[PivotList+1].docid() = PivotID do
12
                 PivotList ← PivotList +1
13
14
             end
             break
15
          end
16
          PivotList ← PivotList +1
      if UpperBound \leq \theta or PivotID = MaxDocID then
                                                                  // No pivot. Exit.
         break
      end
21
      if \mathcal{P}[0].docid() = PivotID then
                                             // Can evaluate the pivot document.
22
          Score ← 0
23
          for List ← 0 to PivotList do
24
             Score \leftarrow Score +\mathcal{P}[List].score()
25
             \mathcal{P}[\text{List}].next()
26
          end
27
          Heap.push((PivotID, Score))
28
          \theta \leftarrow \text{Heap.}min()
29
         SortByDocID(P)
30
31
      end
      else
                                         // Need to align lists with the pivot.
32
          while \mathcal{P}[PivotList] = PivotID do
33
             PivotList ← PivotList -1
34
          end
35
          P[PivotList].nextGEQ(PivotID)
36
         BubbleDown(P, PivotList)
      end
38
39 end
40 return Heap
```

We didn't find a pivot - we are done!

Algorithm 3.3: WAND processing.

```
Input: An array \mathcal{P} of n postings cursors, the largest document identifier, MaxDocID,
            and the number of desired results, k.
  Output: The top-k documents.
 1 Heap ← {}
2 θ ← 0
 3 SortByDocID(P)
 4 while true do
      UpperBound ← 0
      PivotList ← 0
      PivotID ← MaxDocID
      while PivotList < n \text{ do}
                                     // Find the pivot list and pivot document.
          UpperBound \leftarrow UpperBound +\mathcal{P}[PivotList].upperbound()
          if UpperBound > \theta then
10
              PivotID \leftarrow P[PivotList].docid()
11
             while PivotList+1 < n and P[PivotList+1].docid() = PivotID do
12
                  PivotList ← PivotList +1
14
              end
15
              break
          end
16
          PivotList ← PivotList +1
18
      if UpperBound \leq \theta or PivotID = MaxDocID then
                                                                    // No pivot. Exit.
         break
21
      if \mathcal{P}[0].docid() = PivotID then
                                              // Can evaluate the pivot document.
22
23
          Score \leftarrow 0
          for List ← 0 to PivotList do
24
              Score \leftarrow Score +\mathcal{P}[List].score()
25
             \mathcal{P}[\text{List}].next()
26
          end
27
          Heap.push((PivotID, Score))
28
          \theta \leftarrow \text{Heap.}min()
29
          SortByDocID(P)
30
      end
31
      else
                                          // Need to align lists with the pivot.
32
          while \mathcal{P}[PivotList] = PivotID do
33
              PivotList \leftarrow PivotList -1
34
          end
35
          P[PivotList].nextGEQ(PivotID)
36
          BubbleDown(P, PivotList)
      end
38
39 end
40 return Heap
```

If the first list "points to" the pivot document, then all lists "in play" are pointing to the pivot. We score the pivot, try to add it to the heap, and then re-sort the cursors.

Algorithm 3.3: WAND processing.

```
Input: An array \mathcal{P} of n postings cursors, the largest document identifier, MaxDocID,
            and the number of desired results, k.
  Output: The top-k documents.
 1 Heap ← {}
 2 θ ← 0
 3 SortByDocID(P)
 4 while true do
      UpperBound ← 0
      PivotList ← 0
      PivotID ← MaxDocID
      while PivotList < n do
                                    // Find the pivot list and pivot document.
          UpperBound \leftarrow UpperBound +\mathcal{P}[PivotList].upperbound()
          if UpperBound > \theta then
10
              PivotID \leftarrow P[PivotList].docid()
11
              while PivotList+1 < n and P[PivotList+1].docid() = PivotID do
12
                 PivotList ← PivotList +1
14
              end
15
              break
          end
16
          PivotList ← PivotList +1
18
      if UpperBound \leq \theta or PivotID = MaxDocID then
                                                                   // No pivot. Exit.
         break
      end
21
      if \mathcal{P}[0].docid() = PivotID then
22
                                             // Can evaluate the pivot document.
23
          Score \leftarrow 0
          for List ← 0 to PivotList do
24
              Score \leftarrow Score +\mathcal{P}[List].score()
25
             \mathcal{P}[\text{List}].next()
26
          end
27
          Heap.push((PivotID, Score))
28
          \theta \leftarrow \text{Heap.}min()
29
          SortByDocID(P)
30
31
      end
      else
                                          // Need to align lists with the pivot.
32
          while \mathcal{P}[PivotList] = PivotID do
33
              PivotList \leftarrow PivotList -1
34
35
          end
          P[PivotList].nextGEQ(PivotID)
36
          BubbleDown(P, PivotList)
38
39 end
40 return Heap
```

Otherwise, there are lists that point to documents smaller than the pivot. We need to move them up to the pivot before we score. Note that this may require partial sorting (BubbleDown)

MaxScore

Algorithm 3.2: DAAT MaxScore processing.

```
Input: An array \mathcal{P} of n postings cursors which are sorted increasing on their
             upper-bound values, the largest document identifier, MaxDocID, and the
             number of desired results, k.
   Output: The top-k documents.
 1 Heap ← {}
2 θ ← 0
3 CumulativeBounds ← {}
4 CumulativeBounds[0] \leftarrow \mathcal{P}[0].upperbound()
5 for i \leftarrow 1 to n-1 do
6 | CumulativeBounds[i] \leftarrow CumulativeBounds[i-1] +\mathcal{P}[i].upperbound()
7 end
8 PivotList ← 0
9 PivotID \leftarrow MinimumDocID(\mathcal{P})
10 while PivotID \leq MaxDocID and PivotList < n do
       Score ← 0
       NextCandidate ← MaxDocID
12
       for i \leftarrow PivotList to n-1 do
                                                              // Score essential lists.
          if \mathcal{P}[i].docid() = PivotID then
               Score \leftarrow Score +\mathcal{P}[i].score()
              P[i].next()
16
           end
17
          if \mathcal{P}[i].docid() < NextCandidate then
18
               NextCandidate \leftarrow P[i].docid()
19
20
           end
21
       end
       for i \leftarrow PivotList - 1 to 0 do // Complete scoring on non-essential lists.
          if Score + CumulativeBounds[i] \leq \theta then
23
               break
24
           end
25
           \mathcal{P}[i].nextGEQ(PivotID)
26
           if \mathcal{P}[i].docid () = PivotID then
               Score \leftarrow Score +\mathcal{P}[i].score()
          end
29
30
       end
       CurrentBound \leftarrow \theta
       Heap.push((PivotID, Score))
       \theta \leftarrow \mathsf{Heap}.min()
       if CurrentBound < \theta then
                                                    // The heap threshold increased.
          while PivotList < n and CumulativeBounds[PivotList] \le \theta do
               PivotList ← PivotList +1
37
          end
       PivotID \leftarrow NextCandidate
40 end
41 return Heap
```

MaxScore

Algorithm 3.2: DAAT MaxScore processing.

```
Input : An array \mathcal{P} of n postings cursors which are sorted increasing on their
             upper-bound values, the largest document identifier, MaxDocID, and the
             number of desired results, k.
  Output: The top-k documents.
1 Heap ← {}
2 0 . 0
3 CumulativeBounds ← {}
4 CumulativeBounds[0] \leftarrow \mathcal{P}[0].upperbound()
5 for i \leftarrow 1 to n-1 do
6 | CumulativeBounds[i] \leftarrow CumulativeBounds[i-1] +\mathcal{P}[i].upperbound()
7 end
8 PivotList ← 0
9 Distall Minimum Desilo(D)
10 while PivotID \leq MaxDocID and PivotList < n do
       Score ← 0
       NextCandidate \leftarrow MaxDocID
12
       for i \leftarrow PivotList to n-1 do
                                                              // Score essential lists.
          if \mathcal{P}[i].docid() = PivotID then
               Score \leftarrow Score +\mathcal{P}[i].score()
              P[i].next()
16
          end
17
          if \mathcal{P}[i].docid() < NextCandidate then
18
               NextCandidate \leftarrow \mathcal{P}[i].docid()
20
          end
21
       end
      for i \leftarrow PivotList - 1 to 0 do // Complete scoring on non-essential lists.
          if Score + CumulativeBounds[i] \leq \theta then
              break
24
          end
25
          \mathcal{P}[i].nextGEQ(PivotID)
26
          if \mathcal{P}[i].docid () = PivotID then
               Score \leftarrow Score +\mathcal{P}[i].score()
          end
29
30
       end
       CurrentBound \leftarrow \theta
       Heap.push((PivotID, Score))
       \theta \leftarrow \text{Heap.}min()
       if CurrentBound < \theta then
                                                    // The heap threshold increased.
          while PivotList < n and CumulativeBounds[PivotList] \le \theta do
               PivotList ← PivotList +1
37
          end
       PivotID ← NextCandidate
40 end
41 return Heap
```

Compute the cumulative sum of the upper-bounds

MaxScore

Algorithm 3.2: DAAT MaxScore processing.

```
Input: An array \mathcal{P} of n postings cursors which are sorted increasing on their
             upper-bound values, the largest document identifier, MaxDocID, and the
             number of desired results, k.
  Output: The top-k documents.
1 Heap ← {}
2 θ ← 0
3 CumulativeBounds ← {}
4 CumulativeBounds[0] \leftarrow \mathcal{P}[0].upperbound()
5 for i \leftarrow 1 to n-1 do
6 | CumulativeBounds[i] \leftarrow CumulativeBounds[i-1] +\mathcal{P}[i].upperbound()
7 end
  Pivotl ist \leftarrow 0
  PivotID \leftarrow MinimumDocID(P)
while PivotID \leq MaxDocID and PivotList < n do
       Score ← 0
       NextCandidate \leftarrow MaxDocID
       for i \leftarrow PivotList to n-1 do
                                                              // Score essential lists
          if \mathcal{P}[i].docid() = PivotID then
               Score \leftarrow Score + P[i].score()
              P[i].next()
          end
          if \mathcal{P}[i].docid() < NextCandidate then
              NextCandidate \leftarrow P[i].docid()
          end
2
      for i — Tivothist — I to 0 do // Complete scotting on non-essential lists
          if Score + CumulativeBounds[i] \leq \theta then
              break
24
          end
25
          \mathcal{P}[i].nextGEQ(PivotID)
26
          if \mathcal{P}[i].docid () = PivotID then
              Score \leftarrow Score +\mathcal{P}[i].score()
          end
29
30
       end
       CurrentBound \leftarrow \theta
       Heap.push((PivotID, Score))
       \theta \leftarrow \mathsf{Heap}.min()
       if CurrentBound < \theta then
                                                    // The heap threshold increased.
          while PivotList < n and CumulativeBounds[PivotList] \le \theta do
              PivotList ← PivotList +1
37
          end
       PivotID ← NextCandidate
40 end
41 return Heap
```

The pivot document is the minimum document in the essential lists. We score this document, and also track the next pivot candidate.

MaxScore

Algorithm 3.2: DAAT MaxScore processing.

```
Input : An array \mathcal{P} of n postings cursors which are sorted increasing on their
             upper-bound values, the largest document identifier, MaxDocID, and the
             number of desired results, k.
  Output: The top-k documents.
1 Heap ← {}
2 θ ← 0
3 CumulativeBounds ← {}
4 CumulativeBounds[0] \leftarrow \mathcal{P}[0].upperbound()
5 for i \leftarrow 1 to n-1 do
6 | CumulativeBounds[i] \leftarrow CumulativeBounds[i-1] +\mathcal{P}[i].upperbound()
7 end
8 PivotList ← 0
9 PivotID \leftarrow MinimumDocID(P)
10 while PivotID \leq MaxDocID and PivotList < n do
       Score ← 0
       NextCandidate \leftarrow MaxDocID
12
       for i \leftarrow PivotList to n-1 do
                                                              // Score essential lists.
          if \mathcal{P}[i].docid() = PivotID then
               Score \leftarrow Score +\mathcal{P}[i].score()
              P[i].next()
16
          end
17
          if \mathcal{P}[i].docid() < NextCandidate then
               NextCandidate \leftarrow P[i].docid()
20
          end
21
      for i \leftarrow PivotList - 1 to 0 do // Complete scoring on non-essential lists.
          if Score + CumulativeBounds[i] \leq \theta then
23
               break
24
25
          end
          \mathcal{P}[i].nextGEQ(PivotID)
26
           if \mathcal{P}[i].docid() = PivotID then
               Score \leftarrow Score +\mathcal{P}[i].score()
28
          end
29
       end
30
       CurrentBound \leftarrow \theta
       Heap.push((PivotID, Score))
       \theta \leftarrow \mathsf{Heap}.min()
       if CurrentBound < \theta then
                                                     // The heap threshold increased.
          while PivotList < n and CumulativeBounds[PivotList] \le \theta do
               PivotList ← PivotList +1
37
          end
       PivotID ← NextCandidate
40 end
41 return Heap
```

The pivot document may also occur in the non-essential lists; we complete scoring the pivot here.

MaxScore

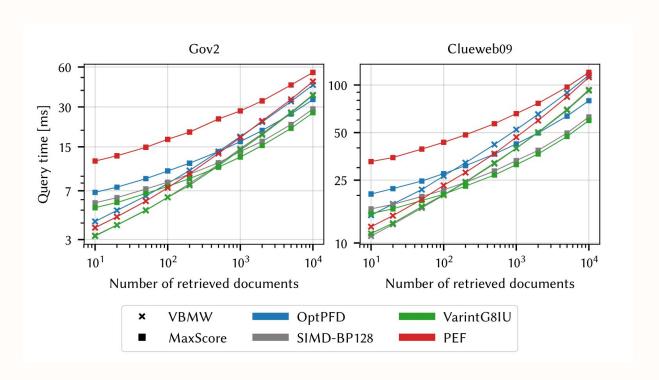
Algorithm 3.2: DAAT MaxScore processing.

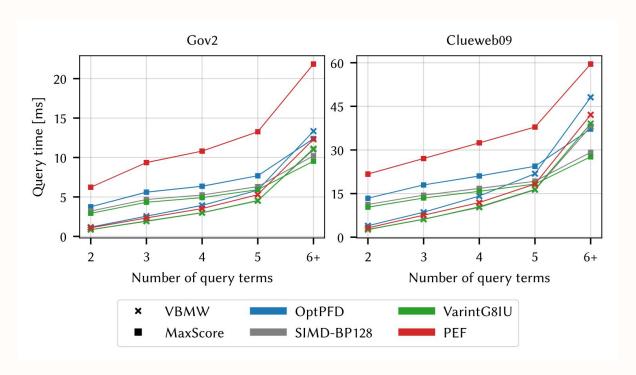
```
Input: An array \mathcal{P} of n postings cursors which are sorted increasing on their
            upper-bound values, the largest document identifier, MaxDocID, and the
             number of desired results, k.
  Output: The top-k documents.
1 Heap ← {}
2 θ ← 0
3 CumulativeBounds ← {}
4 CumulativeBounds[0] \leftarrow \mathcal{P}[0].upperbound()
5 for i \leftarrow 1 to n-1 do
6 | CumulativeBounds[i] \leftarrow CumulativeBounds[i-1] +\mathcal{P}[i].upperbound()
7 end
8 Pivotl ist ← 0
9 PivotID \leftarrow MinimumDocID(P)
10 while PivotID \leq MaxDocID and PivotList < n do
       Score ← 0
       NextCandidate \leftarrow MaxDocID
       for i \leftarrow PivotList to n-1 do
                                                              // Score essential lists.
          if \mathcal{P}[i].docid() = PivotID then
               Score \leftarrow Score +\mathcal{P}[i].score()
              P[i].next()
16
          end
17
          if \mathcal{P}[i].docid() < NextCandidate then
               NextCandidate \leftarrow P[i].docid()
20
          end
21
       end
      for i \leftarrow PivotList -1 to 0 do // Complete scoring on non-essential lists.
          if Score + CumulativeBounds[i] \leq \theta then
               break
24
          end
25
          \mathcal{P}[i].nextGEQ(PivotID)
26
          if \mathcal{P}[i].docid () = PivotID then
               Score \leftarrow Score +\mathcal{P}[i].score()
          end
29
       end
30
       CurrentBound \leftarrow \theta
       Heap.push((PivotID, Score))
       \theta \leftarrow \mathsf{Heap}.min()
       if CurrentBound < \theta then
                                                    // The heap threshold increased.
          while PivotList < n and CumulativeBounds[PivotList] \le \theta do
               PivotList ← PivotList +1
37
          end
38
       PivotID ← NextCandidate
40 end
41 return Heap
```

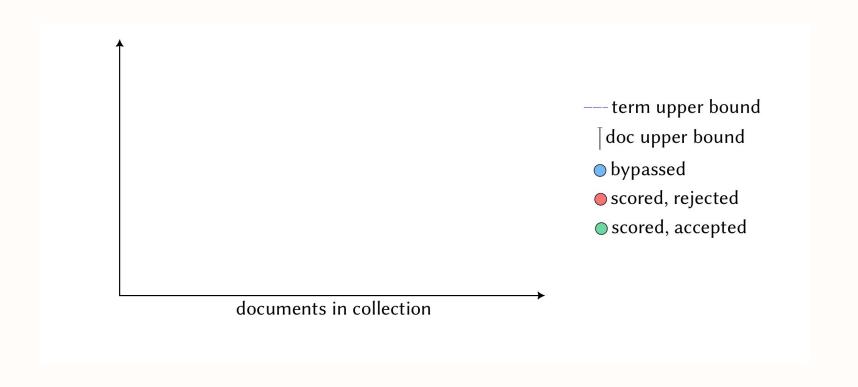
Now we check to see if the pivot can enter the heap. If so, we may need to adjust the boundary between the essential and non-essential lists.

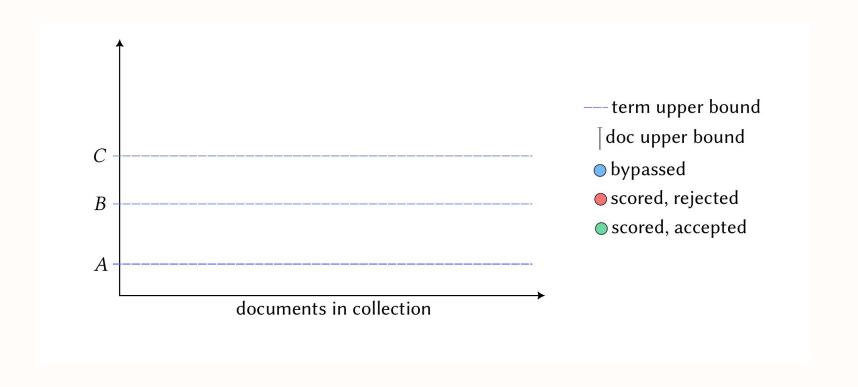
- WAND typically performs well for short queries, and small values of k
 - As query length increases, the sorting operations become expensive.
 - As k increases, dynamic pruning becomes less effective, as the heap threshold is easier to beat, meaning we score more documents.

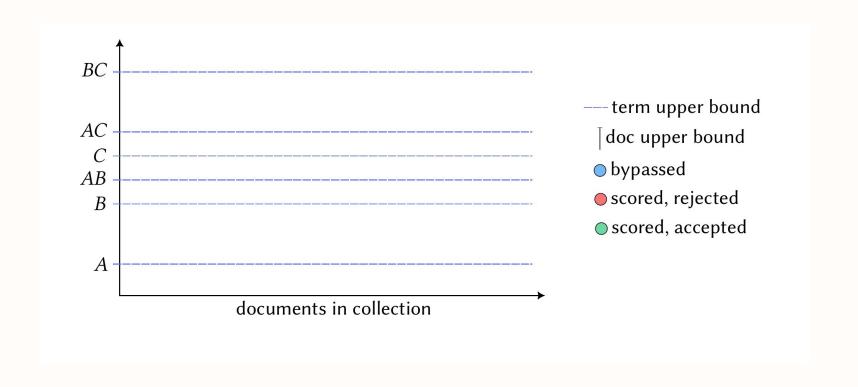
- MaxScore typically performs well for long(er) queries, and large(r) values of k
 - No sorting required during processing!

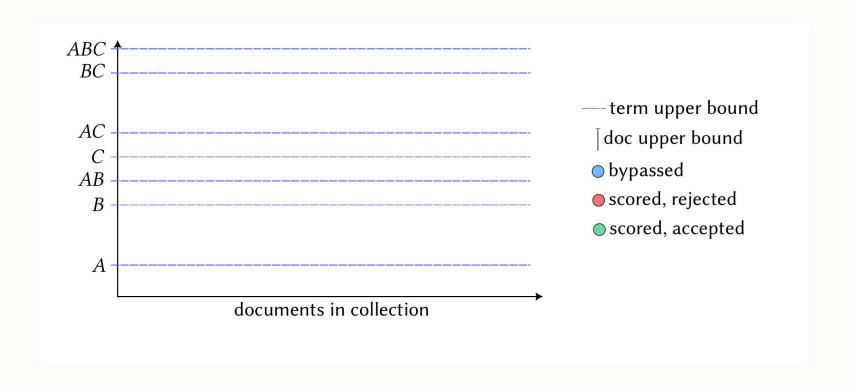


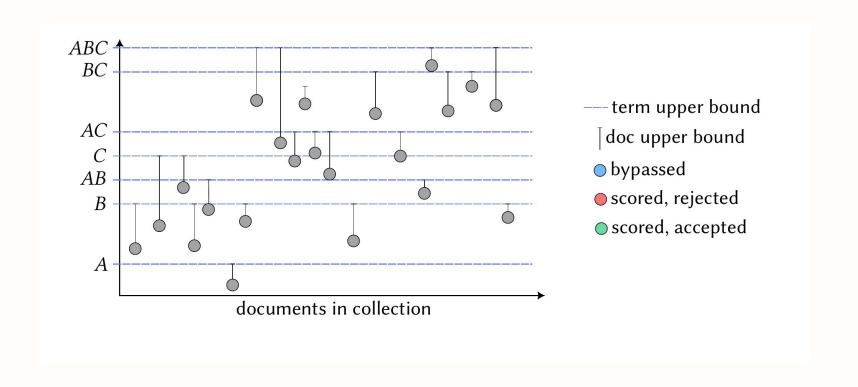


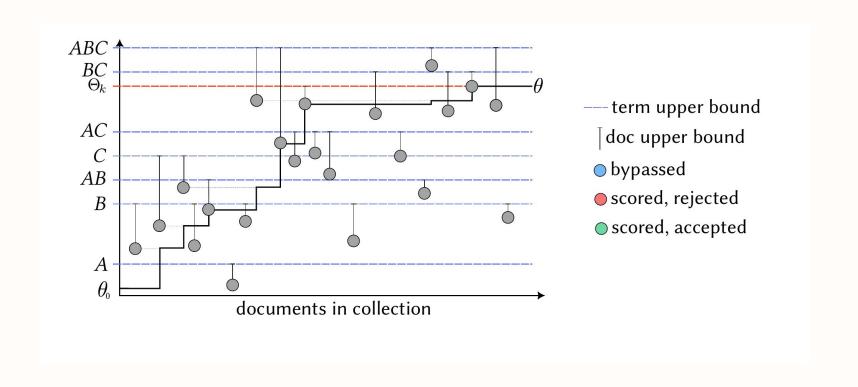


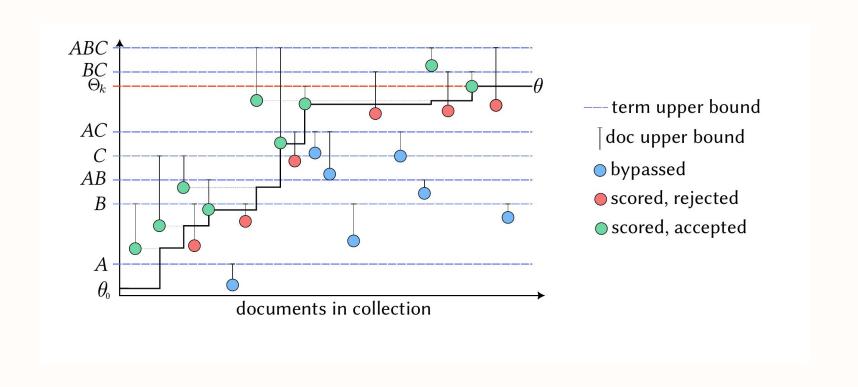












Upper-Bound Estimation

MaxScore and WAND use list-wise upper-bounds to make estimations on document scores. These estimations are used for "go" or "no go" choices on document scoring.

But using the maximum list-wise score may not be a good estimate. How can we do better?

One approach is to store a *per-block* upper-bound in addition to the list-wise upper-bound score.

The list-wise upper-bounds drive the initial selection of a candidate; then a localized upper-bound allows for a more accurate decision to be made before proceeding.

The obvious **downside** is the additional space consumption. But this is typically small, and these bounds can be compressed.

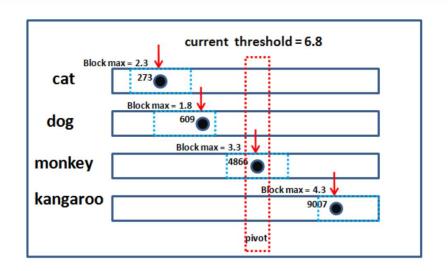


Figure 4: An example showing why directly using block max scores does not work.

We **cannot** simply use the block-max scores to decide which pivot to score, or we may skip documents that should be in the top-k.

Some optimisations intentionally do this, resulting in *unsafe* retrieval.

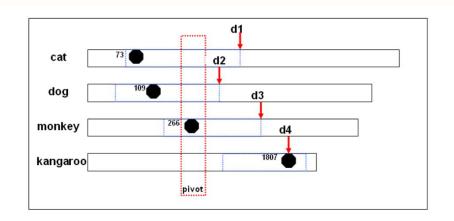


Figure 5: An example showing how GetNewCandidate() works. Assume 266 is the pivot and it fails to make it into the top results. In this case, we enable better skipping by choosing $\min(d1,d2,d3,d4)$ as the next possible candidate, instead of 266+1

However, we can use the block-max bounds to make better decisions on what to process next!

In this case, the current block "configuration" cannot yield a document that will be admitted into the top-k.

Many versions of Block-Max MaxScore and Block-Max WAND

- Window-Based Blocks
- Live-Block Pruning
- Conditional Skipping
- Hybrid Approaches (LazyBM)
- Many more...

Intuitively, all of these algorithms are variations that improve the plain BMM/BMW algorithms through specific observations; the literature is dense!

Variable-Sized Blocks

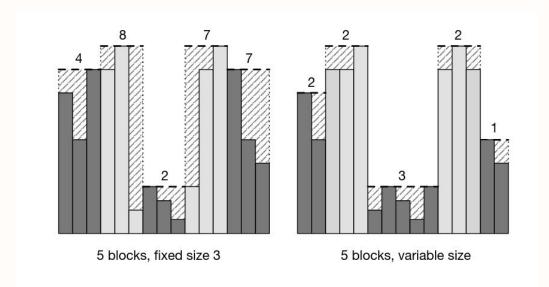


Figure 1: Block errors in constant (left) and variable (right) block partitioning.

Intuition: Fixed-size blocks may cause large within-block errors. Instead, find a variable-length partition that reduces error rate.

The Ranker does Matter!

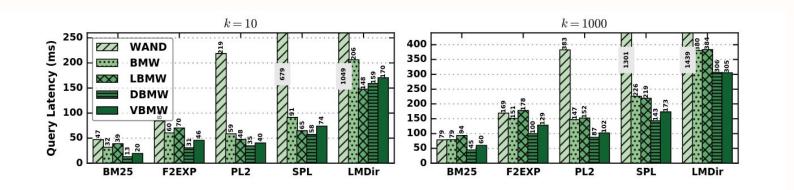


Figure 2: Mean query latency (in milliseconds) of five WAND-based strategies across five ranking models. Latency is reported for top-k queries for k = 10 and k = 1000.

Finding the Best of Both Worlds: Faster and More Robust Top-k Document Retrieval - Khattab, Hammoud, Elsayed - SIGIR 2020

The Ranker does Matter!

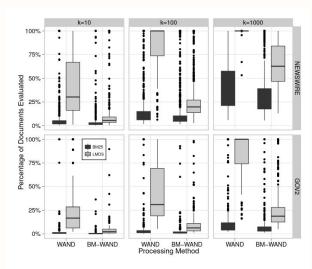


Figure 2: Number of documents scored, expressed as a fraction of the total number of documents containing any of the terms in that query. Three different retrieval depths are tested, for two different collections, for two different similarity computations, and for both WAND and BM-WAND processing.

Exploring the Magic of WAND - Petri, Culpepper, Moffat, ADCS 2013.

The Ranker does Matter!

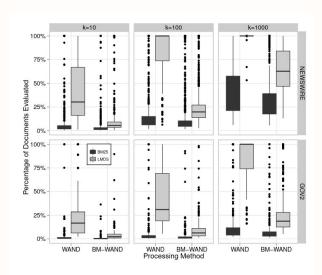


Figure 2: Number of documents scored, expressed as a fraction of the total number of documents containing any of the terms in that query. Three different retrieval depths are tested, for two different collections, for two different similarity computations, and for both WAND and BM-WAND processing.

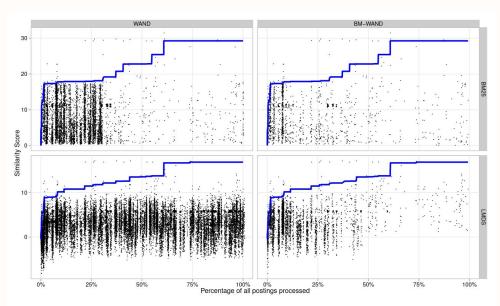
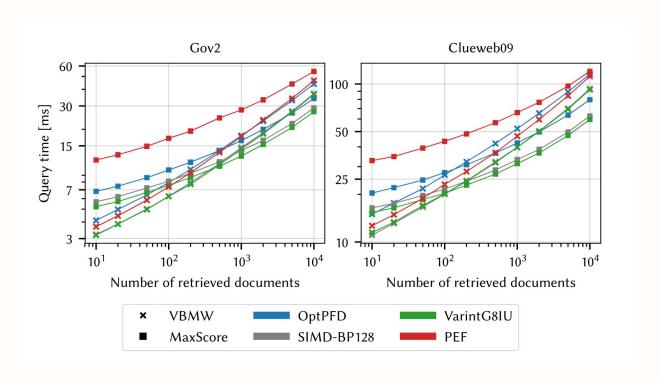
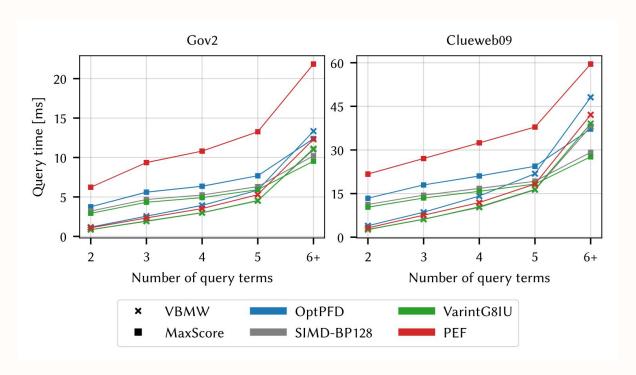


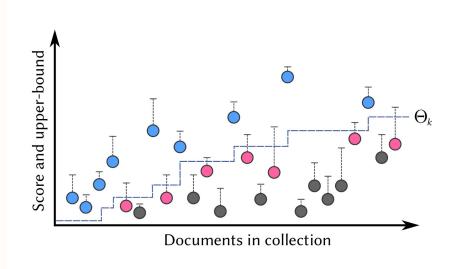
Figure 3: Distribution of evaluated documents as the postings lists for query "north korean counterfeiting" (topic 808) are evaluated for both WAND and BM-WAND processing using LMDS and BM25 similarity measures.

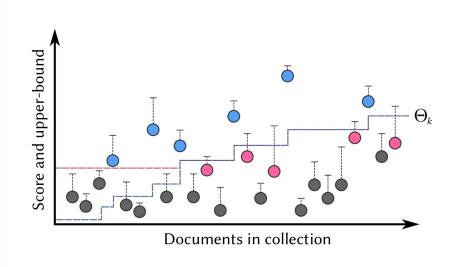
MaxScore vs VBMW

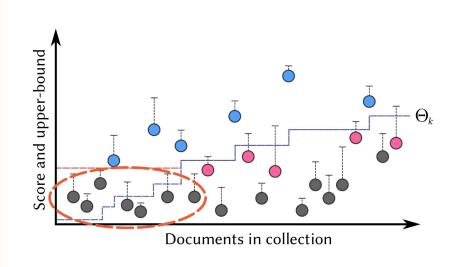


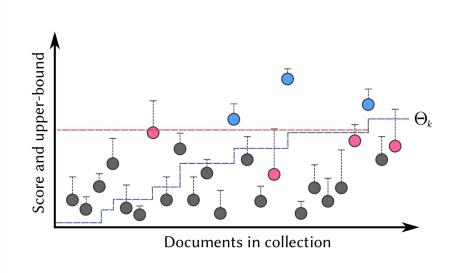
MaxScore vs VBMW

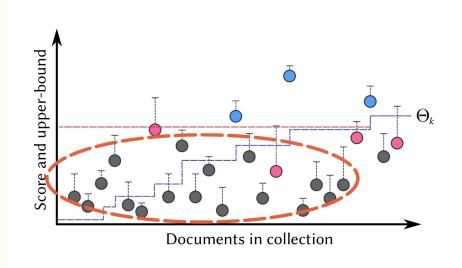










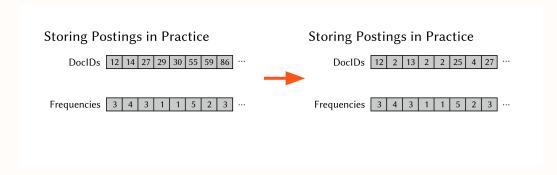


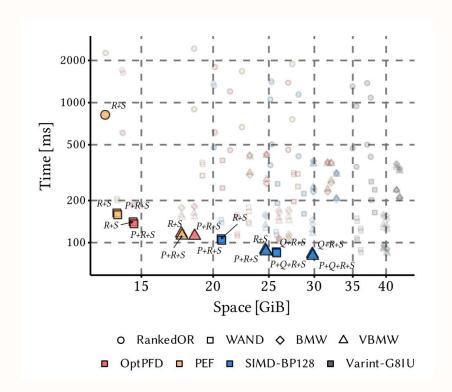
Index Reordering (Document Identifier Reassignment)
 If we can "cluster" similar documents together, we will
 get lots of runs of "1"s in our postings lists

Recall: Postings are strictly increasing on document id, so we *delta* code them.

Index Reordering (Document Identifier Reassignment)
 If we can "cluster" similar documents together, we will
 get lots of runs of "1"s in our postings lists

Recall: Postings are strictly increasing on document id, so we *delta* code them.





Experiment: All combinations of **p**riming, **q**uantization, **r**eordering, **s**topping on Clueweb12B (52 million web documents).

Efficiency innovations are, broadly speaking, additive!

Optimizations	RankedOR	WAND	BMW	VBMW
None	135.7	49.1	45.6	45.6
One	87.9 (×1.5)	$26.6 (\times 1.8)$	19.2 (\times 2.4)	17.9 $(\times 2.6)$
Two	$62.6 (\times 2.2)$	$20.3 (\times 2.4)$	14.8 (×3.1)	$13.3 (\times 3.4)$
Three	$56.2 (\times 2.4)$	15.5 $(\times 3.2)$	12.8 (×3.6)	$10.9 (\times 4.2)$
All	54.7 (×2.5)	13.8 (×3.6)	11.6 (×3.9)	9.7 (×4.7)

Between $2.5 \times$ to $4.7 \times$ speedups.

Taking stock of the current state of top-k query processing: Can retrieve the top 1000 candidates on a 52 million document collection in < 10 milliseconds.

Experiment: All combinations of **p**riming, **q**uantization, **r**eordering, **s**topping on Clueweb12B (52 million web documents).

Efficiency innovations are, broadly speaking, additive!

Rule(s) of Thumb

- Use MaxScore when queries are long, or k is large; Use VBMW otherwise;
- 2. **Stopping** is almost always a good idea;
- 3. Use **index reordering** if you can afford it (offline cost);
- 4. If you don't know which (statistical) ranker to use, just stick to BM25 it is fast, and well behaved;
- 5. Empirical experimentation is always beneficial!

Session I: Indexing and Retrieval

Practice

Indexing and Querying with PISA

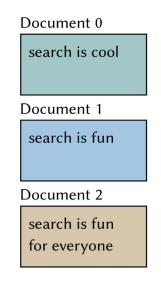
We will now work through **Section 1** of the practical.

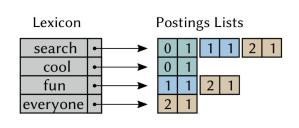
Tutorial: https://shorturl.at/VExpG

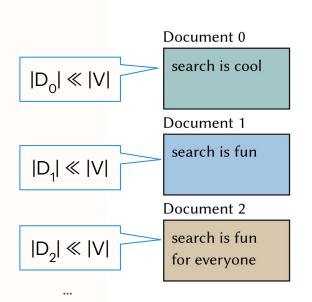
Aka: https://github.com/pisa-engine/pisa/blob/main/test/docker/tutorial/instructions.md

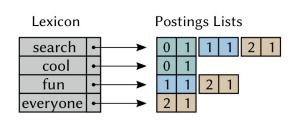
Session 1.5: Discussion & Coffee

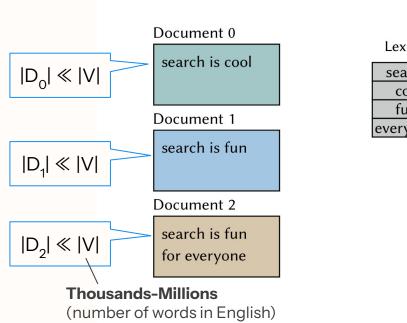
Session II: Learned Sparse Retrieval

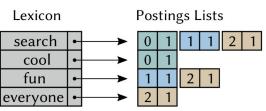


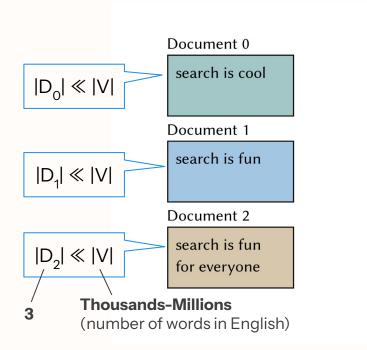


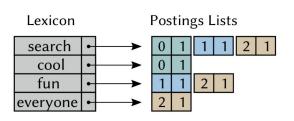


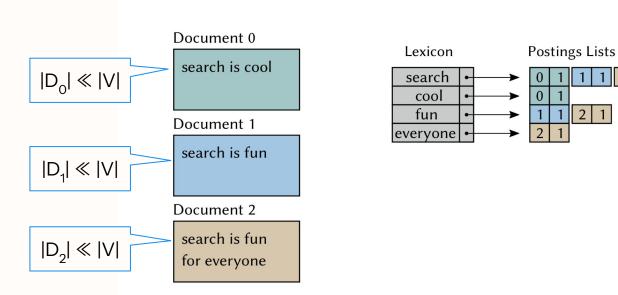




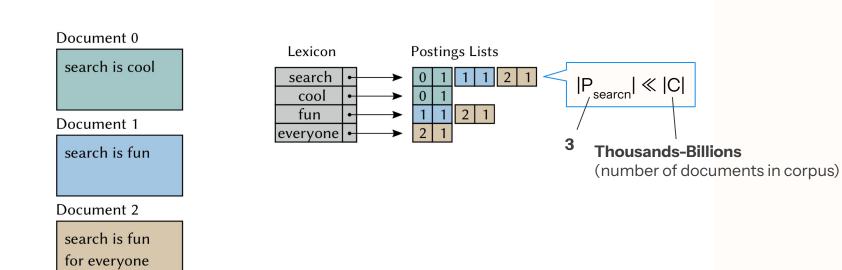


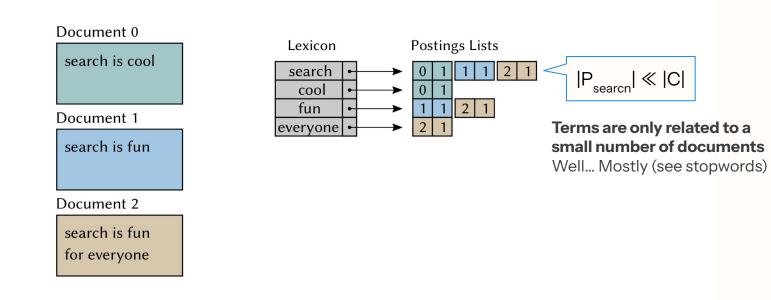


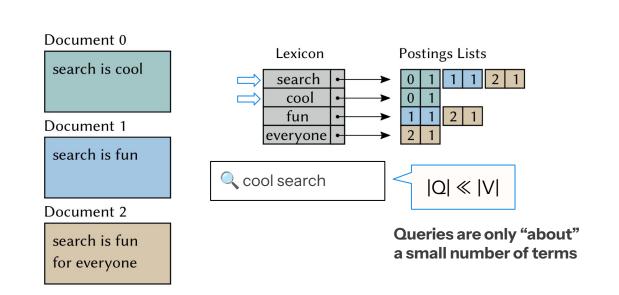




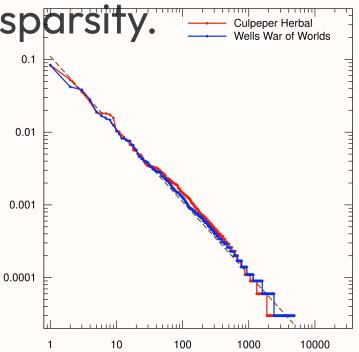
Documents are only "about" a small number of terms







Up to now, we've relied on the presence of terms in documents/queries to enforce



This works because of **Zipf's Law** – very few terms typically have high frequently in a corpus.

A plot of the frequency of each word as a function of its frequency rank for two English language texts. CC BY-SA 4.0

Working Example

Query: F1 winner

Document: Max Verstappen says 3rd Formula One world championship title is his 'best one' so far

LOSAIL, Qatar (AP) — Max Verstappen believes his third Formula One title is his best yet.

Clinching the championship in a sprint race Saturday in Qatar didn't pack the emotional impact of his dramatic, controversial last-lap overtake of Lewis Hamilton for the 2021 title. Still, the Red Bull driver thinks his relentlessly consistent 2023 season has been his greatest so far.

"This one is the best one," Verstappen said. "I think the first one was the most emotional one because that's when your dreams are fulfilled in Formula One. But this one definitely in my opinion has been my best year also for consecutive wins and stuff. The car itself has been probably in the best shape as well. This one is probably (the one) I'm most proud of in a way because of consistency."

. . .

Source: AP Online News

Maps a query and document to a sparse "bag-of-words" representation using:

- **TF** (importance of term to the document, based on repetition)
- **IDF** (relative importance of term to the query, based on how many documents it appears in)
- Other Lexical Signals (e.g., document length for normalization, etc.)

$$\sum_{t \in q} \log \left(\frac{N}{df_t} \right) \cdot \frac{(k_1 + 1) \cdot t f_{td}}{k_1 \cdot \left(1 - b + b \cdot \left(\frac{L_d}{L_{avg}} \right) \right) + t f_{td}}$$

Maps a query and document to a sparse "bag-of-words" representation using:

- **TF** (importance of term to the document, based on repetition)
- **IDF** (relative importance of term to the query, based on how many documents it appears in)
- Other Lexical Signals (e.g., document length for normalization, etc.)

Problems with term-based representations:

(1) TF and IDF aren't always good estimators of term importance

Example: "Saturday" probably isn't more important than "Sunday" in the article

(if anything, "Sunday" is more important to the document – that's when the race was)

Problems with term-based representations:

(1) TF and IDF aren't always good estimators of term importance

Example: "Saturday" probably isn't more important than "Sunday" in the article (if anything, "Sunday" is more important to the document – that's when the race was)

(2) Lexical mismatch: variations of terms that appear in the query/document

Example: "F1" is very important in the article, but isn't mentioned (similar terms are, though, e.g., "Formula One")

Learned Sparse Retrieval (LSR) uses neural networks to produce better bag-of-words representations

(addressing the aforementioned issues)

Learned Sparse Retrieval

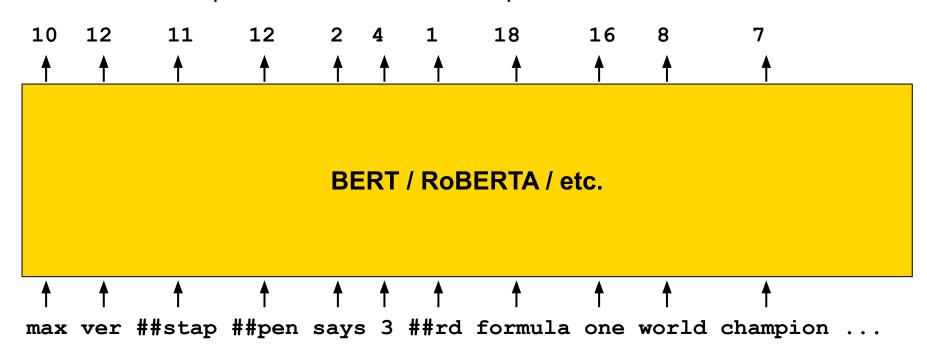
Three key techniques:

- (1) Term weighting (addresses importance estimation)
- (2) Expansion (addresses lexical mismatch)
- (3) Sparsification (manages compute & storage efficiency of (2))

(1) Term Weighting

Main idea: Use a neural network to estimate the "importance" of each token.

Implemented as a token-level prediction task.



(1) Term Weighting

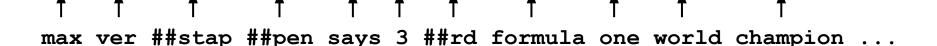
Main idea: Use a neural network to estimate the "importance" of each token.

```
{formula: 18, one: 16, ver: 12, ##pen: 12, ##stap: 11, ...}

10 12 11 12 2 4 1 18 16 8 7

↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑
```

BERT / RoBERTA / etc.



(1) Term Weighting

Term weighting can be applied to the document [1] and/or the query [2].

They can even learn what to remove (weight=0) [3]

[1] Zhuyun Dai, Jamie Callan. **Context-Aware Sentence/Passage Term Importance Estimation** For First Stage Retrieval. arxiv 2019. <u>link</u>

[2] Sean MacAvaney, Franco Maria Nardini, Raffaele Perego, Nicola Tonellotto, Nazli Goharian, Ophir Frieder. **Expansion via Prediction of Importance with Contextualization.** SIGIR 2020. <u>Link</u>

[3] Mackenzie et al. Efficiency Implications of Term Weighting for Passage Retrieval. SIGIR 2020. link

Goal: Identify new terms to add to the document (and estimate their importance)

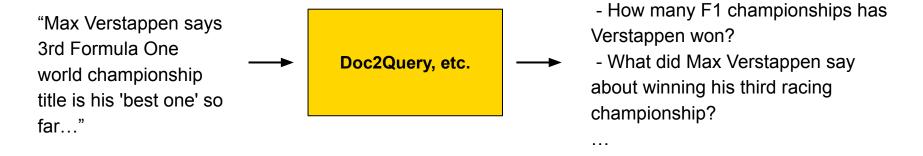
Two main approaches:

(2a) External Expansion

(2b) Masked Language Modeling (MLM) Expansion

(2a) External Expansion

Use another model (such as Doc2Query [1]) to add expansion tokens to document [2], then apply term weighting.

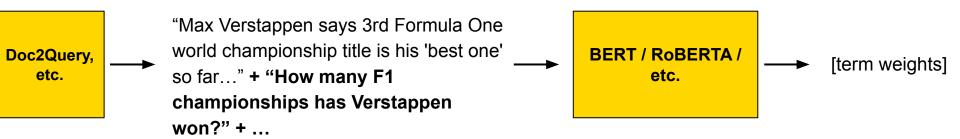


[1] Rodrigo Nogueira, Wei Yang, Jimmy Lin, Kyunghyun Cho. Document Expansion by Query Prediction. arxiv 2019. link

[2] Antonio Mallia, Omar Khattab, Nicola Tonellotto, Torsten Suel. Learning Passage Impacts for Inverted Indexes. SIGIR 2021. link

(2a) External Expansion

Use another model (such as Doc2Query [1]) to add expansion tokens to document [2], then apply term weighting.



[1] Rodrigo Nogueira, Wei Yang, Jimmy Lin, Kyunghyun Cho. Document Expansion by Query Prediction. arxiv 2019. Link

[2] Antonio Mallia, Omar Khattab, Nicola Tonellotto, Torsten Suel. Learning Passage Impacts for Inverted Indexes. SIGIR 2021. link

(2b) Masked Language Modeling Expansion

Use the model's Masked Language Modelling head to get expansion terms.

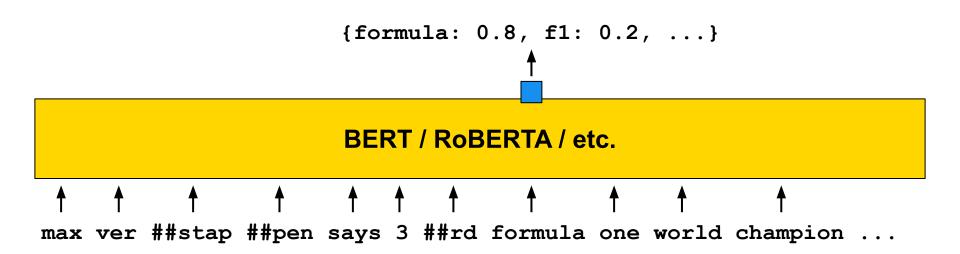
(2b) Masked Language Modeling Expansion

Use the model's Masked Language Modelling head to get expansion terms.

Review: Pre-training a BERT-like language model

(2b) Masked Language Modeling Expansion

Use the model's Masked Language Modelling head to get expansion terms.



(2b) Masked Language Modeling Expansion

Use the model's Masked Language Modelling head to get expansion terms.

{formula: 18, one: 16, ver: 12, ##pen: 12, **f1: 10**, ...}

(2b) MLM Expansion can be learned end-to end, and can be applied to the document [1] and/or the query [2].

- [1] Sean MacAvaney, Franco Maria Nardini, Raffaele Perego, Nicola Tonellotto, Nazli Goharian, Ophir Frieder. **Expansion via Prediction of Importance with Contextualization.** SIGIR 2020. <u>link</u>
- [2] Thibault Formal, Benjamin Piwowarski and Stéphane Clinchant. **SPLADE: Sparse Lexical and Expansion Model for First Stage Ranking.** SIGIR 2021. <u>Link</u>
- [3] Zhuang and Zuccon. Fast Passage Re-ranking with Contextualized Exact Term Matching and Efficient Passage Expansion. <u>link</u>

(3) Sparsification

MLM Expansion actually makes the output vectors dense – it computes an importance score for every token in the lexicon (10's of thousands of dimensions).

We could index these vectors in dense vector stores, but they're very large, and this makes retrieval slow.

Instead, we sparsify, allowing vectors to be used in typical inverted indexes (e.g., Lucene)

Sparsification approaches:

(3a): Top-K Pruning [1] (post-hoc pruning of lowest vector dimensions)

(3b): FLOPS Regularisation [2] (end-to-end optimisation to push dimensions to zero)

(3c): DF-FLOPS Regularisation [3] (objective to reduce number of terms per document)

[1] Sean MacAvaney, Franco Maria Nardini, Raffaele Perego, Nicola Tonellotto, Nazli Goharian, Ophir Frieder. **Expansion via Prediction of Importance with Contextualization.** SIGIR 2020. <u>link</u>

[2] Thibault Formal, Benjamin Piwowarski and Stéphane Clinchant. SPLADE: Sparse Lexical and Expansion Model for First Stage Ranking. SIGIR 2021. Link

[3] Porco et al. An Alternative to FLOPS Regularization to Effectively Productionize SPLADE-doc. SIGIR 2025. link

Which LSR methods produce better relevance estimates?

In general:

Query Weighting (query expansion only helps a little when document is expanded)

+

Document MLM Expansion

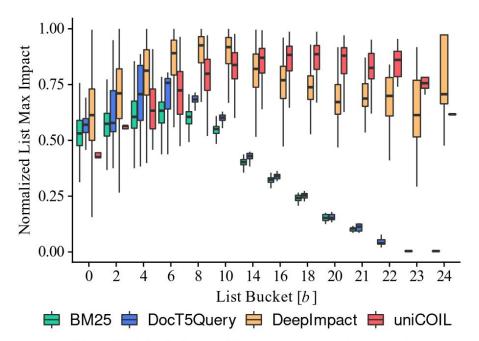
+

Regularisation (for optimal trade-offs) or Top-K (for flexibility)

(Find out more in the Neural Lexical Search with Learned Sparse Retrieval tutorial this afternoon.)

How do LSR methods affect retrieval algorithms?

Term Weighting greatly affects impact distributions



Normalized maximum list impact distribution stratified by list length buckets $b \in [2^b, 2^b+1)$.

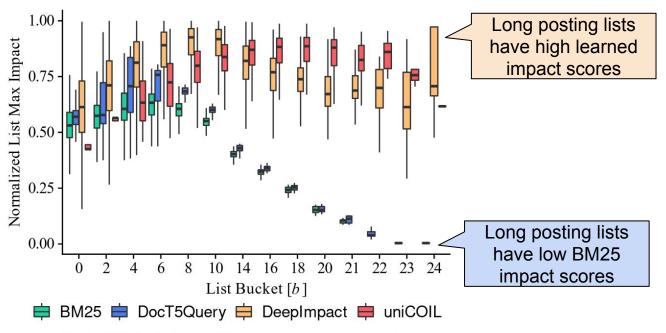
Wacky Weights in Learned Sparse Representations and the Revenge of Score-at-a-Time Query Evaluation

Joel Mackenzie, Andrew Trotman, Jimmy Lin

¹ School of Computing and Information Systems, The University of Melbourne, Australia ² Department of Computer Science, University of Otago, Dunedin, New Zealand

³ David R. Cheriton School of Computer Science, University of Waterloo, Ontario, Canada

Term Weighting greatly affects impact distributions



Wacky Weights in Learned Sparse Representations and the Revenge of Score-at-a-Time Query Evaluation

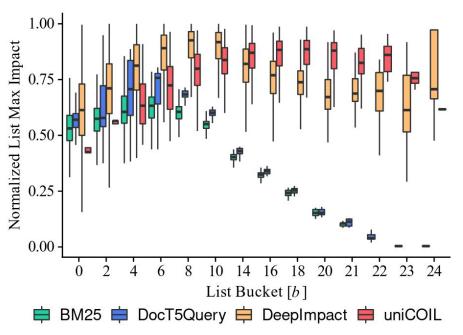
Joel Mackenzie, Andrew Trotman, Jimmy Lin

3

¹ School of Computing and Information Systems, The University of Melbourne, Australia ² Department of Computer Science, University of Otago, Dunedin, New Zealand

³ David R. Cheriton School of Computer Science, University of Waterloo, Ontario, Canada

Term Weighting greatly affects impact distributions



! This means we need to visit more blocks in long posting lists, increasing traversal time

Wacky Weights in Learned Sparse Representations and the Revenge of Score-at-a-Time Query Evaluation

Joel Mackenzie, Andrew Trotman, Jimmy Lin

¹ School of Computing and Information Systems, The University of Melbourne, Australia ² Department of Computer Science, University of Otago, Dunedin, New Zealand

³ David R. Cheriton School of Computer Science, University of Waterloo, Ontario, Canada

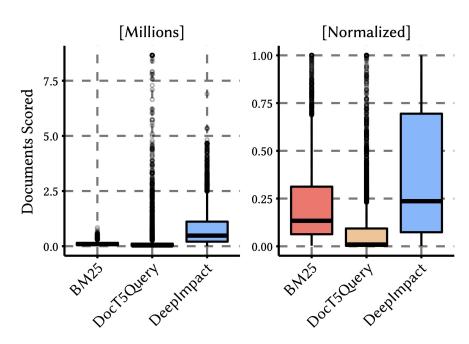
Term Weighting greatly affects impact distributions



Distributions Make or Break Efficiency

Efficient top-k processing algorithms use term upper-bound scores to bypass documents which cannot score highly.

- These algorithms find it more difficult to effectively prune the search space with the impact distributions from learned sparse models.
- More documents are scored; query processing is slower.



There are very long posting lists, too.

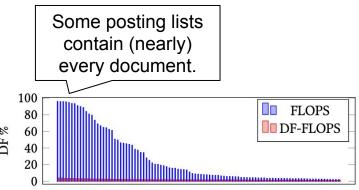


Figure 2: Top 100 DF% in representations produced by SPLADE-Doc with FLOPS and DF-FLOPS on a sample of 100K passages.

There are very long posting lists, too. Why?

Some tokens are repurposed and included in many queries and documents.

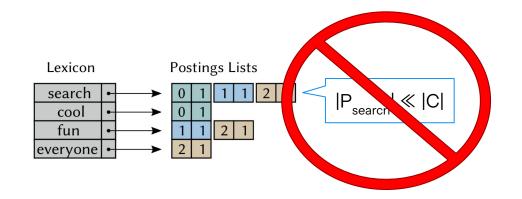
Ends up being partially a pseudo "dense" vector.

```
androgen receptor define →
```

```
('##rogen', 251) ('receptor', 242) ('and', 225) ('receptors', 189)
('hormone', 179) ('definition', 162) ('meaning', 99) ('genus', 89)
('is', 70) (',', 68) 'define', 59 ('the', 56) 'drug', 53) ('for', 46)
('ring', 38) ('gene', 37) ('are', 32) ('god', 25) ('what', 18) ('##rus', 15)
('purpose', 12) ('defined', 10) ('doing', 8 ('a', 4) goal', 4)
```

There are very long posting lists, too.





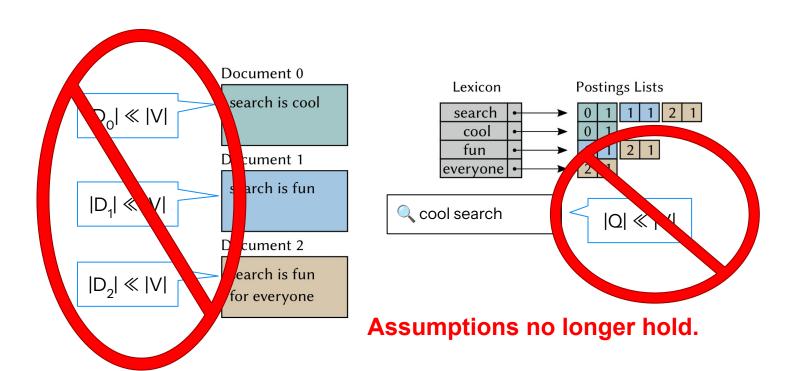
Assumption no longer holds.

Expansion causes much less sparsity in queries/docs

MSMARCO Passages (dev queries)

Method V		Terms in	Documents	Terms in Queries		
	17.1	Total	Unique	Total	Unique	
BM25	2660824	39.8	30.1	5.9	5.8	
BM25-T5	3929111	224.7	51.1	5.9	5.8	
DeepImpact	3514102	4010.0	71.1	4.2	4.2	
uniCOIL-T5	27678	5032.3	66.4	686.3	6.6	
uniCOIL-TILDE	27646	8260.8	107.6	661.1	6.5	
SPLADEv2	28131	10794.8	229.4	2037.8	25.0	

Expansion causes much less sparsity in queries/docs

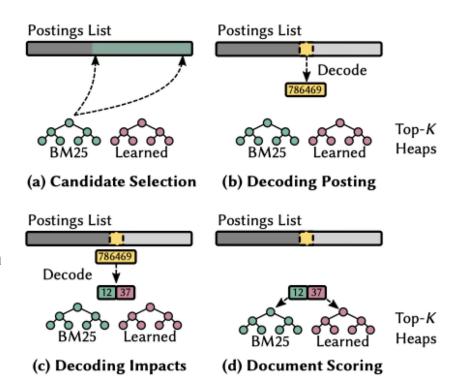


So what can we we do about these problems?

Guided Traversal (GT)

It proposes guided traversal to accelerate top-k processing with learned sparse models.

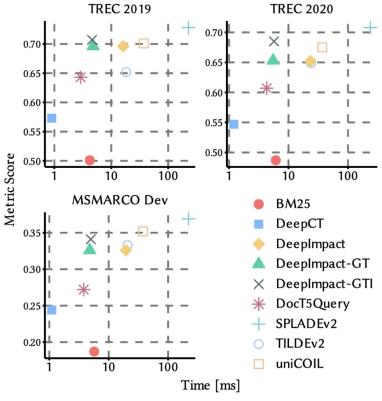
- The original BM25 score for each document (over a DocT5Query expanded index) is stored alongside the learned (DeepImpact) score.
- At query time, BM25 is used to guide the index traversal, but scores are computed via the learned model (Guided Traversal — GT).
- The BM25 score can also be interpolated with the DeepImpact score on-the-fly (Guided Traversal with Interpolation — GTI).



Guided Traversal (GT)

Strategy	MaxScore					
	Mean	Median	an P_{99} RR@2 12.2 0.272 79.6 0.326 15.8 0.326	RR@10		
DocT5Query	3.3	2.8	12.2	0.272		
DeepImpact	19.5	14.0	79.6	0.326		
Dual-Heap	4.6	3.9	15.8	0.326		
Dual-Heap w. Interp. ($\alpha = 0.5$)	4.6	3.9	15.8	0.335		

Guided Traversal (GT)



Mallia et al. "Faster learned sparse retrieval with guided traversal." SIGIR 2022.

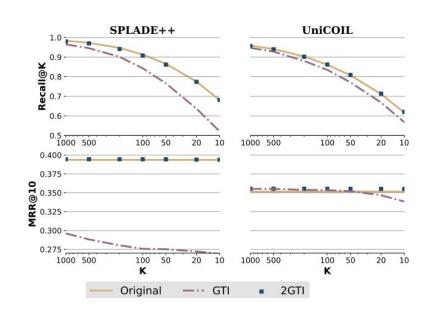
Two-Level Guided Traversal (2GT)

Global pruning – hybrid BM25 + learned upper bounds filter whole posting-list regions.

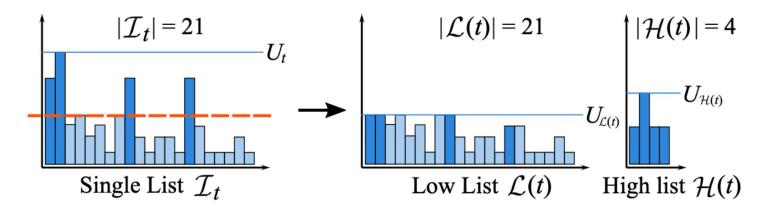
Local pruning – hybrid bounds tighten inside each candidate document.

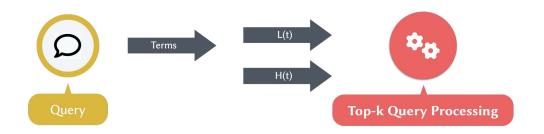
Alignment smoothing makes BM25 weights denser to match the learned index.

Two tunable coefficients (α for global, β for local) cap BM25's influence and prevent over-aggressive skips.



Postings Clipping





Priming can be applied whenever any high-impact list contains k or more postings

$$\theta_0 = \max\{U_{\mathcal{L}(t)} \mid t \in Q \land |\mathcal{H}(t)| \ge k\}$$

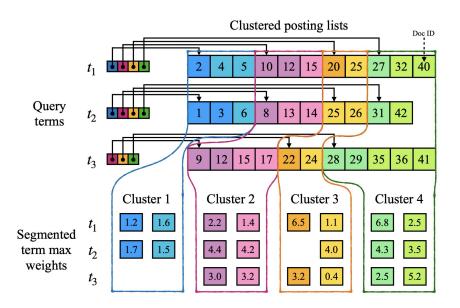
can be used as a priming value for the heap bound, without risking the integrity of the top-k answers.

ASC

Segmented bounds: slice each cluster into n random segments, store per-segment max weights → tighter MaxSBound / AvgSBound

Two-level test: prune cluster if MaxSBound $\leq \theta/\mu$ and AvgSBound $\leq \theta/\eta$; else dive to doc-level pruning at θ/η .

Parameters: $0 < \mu \le \eta \le 1$. Pick μ for aggressiveness, η (often 1) for probabilistic safety.



(a) Cluster skipping index with 2 weight segments per cluster

Qiao et al. Threshold-driven Pruning with Segmented Maximum Term Weights for Approximate Cluster-based Sparse Retrieval. EMNLP 2024.

Seismic

An approximate retrieval solution that trades off exact search for efficiency.

It relies on:

- Concentration of Importance
- Static Document Pruning
- Block Upper Bounds



Figure 1: Fraction of L_1 mass preserved by keeping only the top non-zero entries with the largest absolute value.

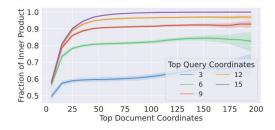
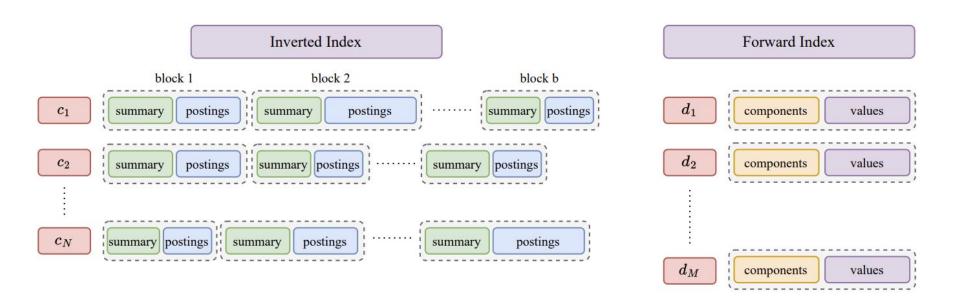


Figure 2: Fraction of inner product (with 95% confidence intervals) preserved by inner product between the top query and document coordinates with the largest absolute value.

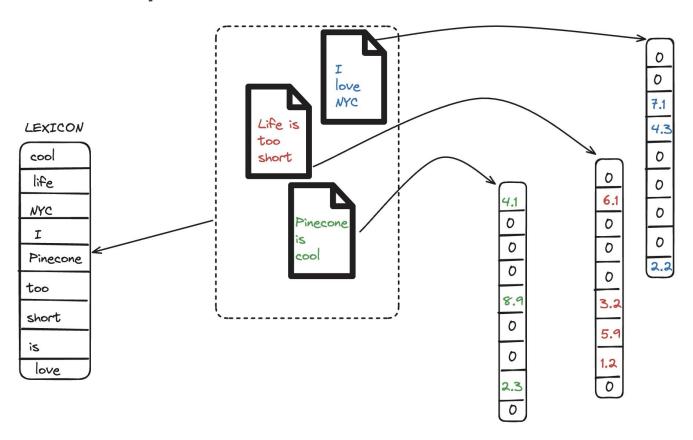
Bruch et al. Efficient Inverted Indexes for Approximate Retrieval over Learned Sparse Representations. SIGIR 2024.

Seismic



Bruch et al. Efficient Inverted Indexes for Approximate Retrieval over Learned Sparse Representations. SIGIR 2024.

Documents as Sparse Vectors



Block-Max Pruning

Vectors of max scores

Block partitioning: The document ID space is split into fixed-size blocks; each block stores a pre-computed vector of its terms' maximum impact scores.

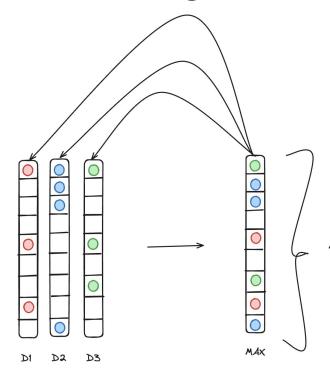
Per-block upper bounds: At query time, term weights are applied to those block-max vectors and summed, producing an overall score upper bound for every block.

Priority-driven evaluation: Blocks are visited in descending order of their upper-bound scores (lazy sorting) until a stopping criterion is satisfied.

Hybrid access: When a block is chosen, lookup switches to a forward-style structure embedded alongside the inverted lists, enabling fast in-block scoring without full postings scans.

Mallia et al. Faster Learned Sparse Retrieval with Block-Max Pruning. SIGIR 2024.

Clustering of Documents

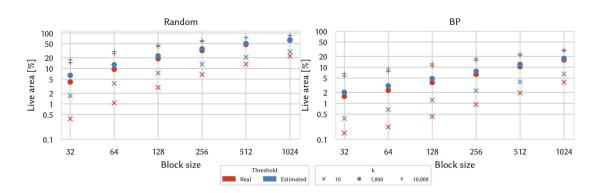


Imagine assigning consecutive docIDs to similar documents.

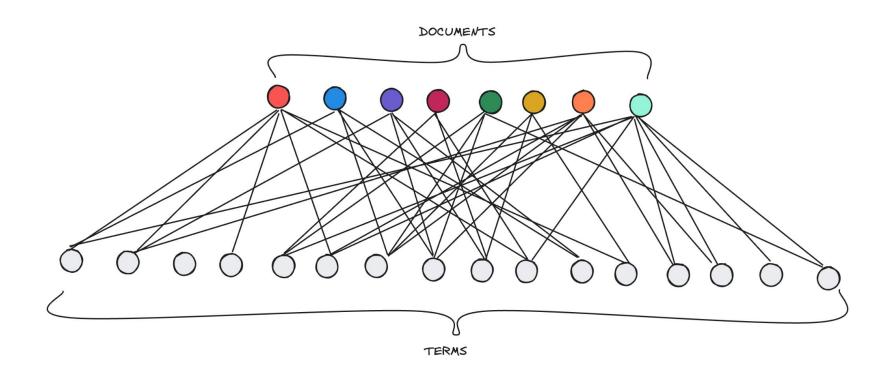
First focus on better compression.

Recently, more focus on faster query processing.

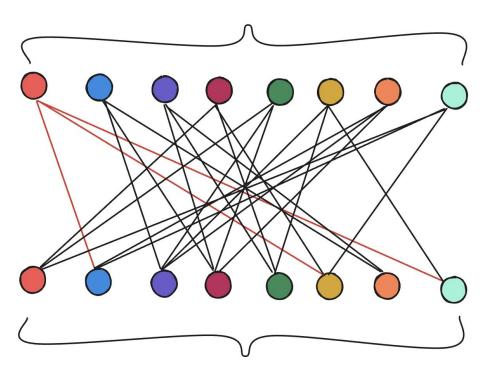




Bipartite Graph Partitioning



Bipartite Graph Partitioning



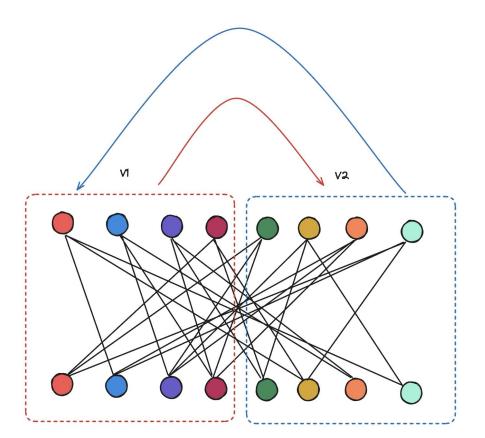
deg(x) is the degree of a node

n is the number of distinct neighbors

The average log gap cost can be proxied with

$$deg(x) \times log(\frac{n}{deg(x) + 1})$$

Graph Bisection



Bisect the graph in two sets

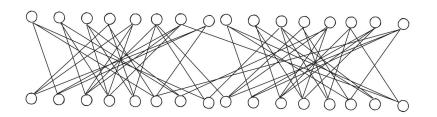
Compute move gains of the vertices. The difference in average logarithmic gap length between remaining and moving

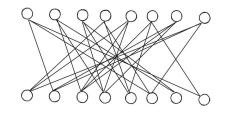
$$deg_1(x) \times log(\frac{n_1}{deg_1(x)+1}) + deg_2(x) \times log(\frac{n_2}{deg_2(x)+1}) = (deg_1(x)-1) \times log(\frac{n_1}{deg_1(x)}) + (deg_2(x)+1) \times log(\frac{n_2}{deg_2(x)+2}) = (deg_1(x)-1) \times log(\frac{n_2}{deg_2(x)}) + (deg_2(x)+1) \times log(\frac{n_2}{deg_2(x)+2}) = (deg_1(x)-1) \times log$$

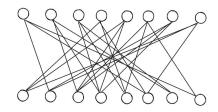
Swap vertices between the two sets

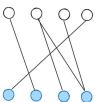
Repeat recursively until stopping condition triggers

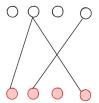
Recursive Graph Bisection

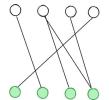


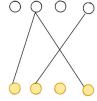




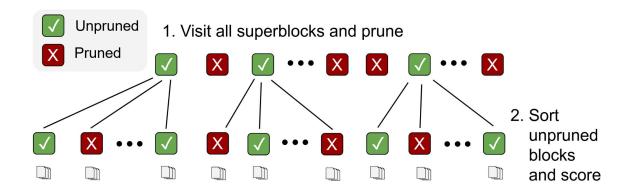








Dynamic Superblock Pruning



Two-tier index: group consecutive document blocks into fixed-size superblocks (e.g., 64 blocks)

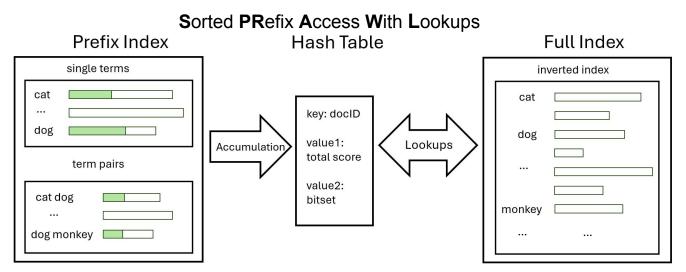
Top-down test: first bound each superblock; prune it if both max- and avg-score $\leq \theta/\mu$, θ/η , then descend to surviving blocks

Dynamic Superblock Pruning

Table 1: Mean response time (ms) and mean reciprocal rank (MRR@10) at a fixed Recall@k budget for SPLADE

Recall	99%		99.5%		99.9%		Rank-Safe		
Budget	MRT	MRR	MRT	MRR	MRT	MRR	MRT	MRR	
k=10									
MaxScore	-	_	_	_	_	-	75.7 (35x)	38.1	
ASC	4.70 (7.5x)	37.9	5.59 (7.8x)	38.1	6.44 (8.2x)	38.1	7.19 (3.3x)	38.1	
Seismic	2.06 (3.3x)	38.1	2.57 (3.6x)	38.2	3.01 (3.8x)	38.4	_	_	
BMP	1.44 (2.3x)	38.1	1.49 (2.1x)	38.1	1.88 (2.4x)	38.2	2.70 (1.3x)	38.1	
SP	0.629	37.7	0.715	37.9	0.785	38.1	2.15	38.1	
k=1000									
MaxScore	_	_	_	_	_	_	124 (12x)	38.1	
ASC	15.8 (9.1x)	38.1	18.9 (9.4x)	38.1	25.4 (5.5x)	38.1	33.5 (3.2x)	38.1	
Seismic	5.72 (3.3x)	38.3	7.18 (3.6x)	38.4	10.5 (2.3x)	38.4	_	_	
BMP	4.99 (2.9x)	38.2	5.25 (2.6x)	38.2	7.26 (1.6x)	38.2	13.9 (1.3x)	38.1	
SP	1.74	37.9	2.01	37.9	4.64	38.2	10.5	38.1	

SPRAWL



It uses a two-tier in-memory index:

- A prefix index that stores top-scoring postings (sorted by impact scores in descending order) for frequent terms and pairs of terms.
- A standard inverted index that containing all standard (single-term) postings and that efficiently supports random lookups.

Gou et al. "Fast and Effective Early Termination for Simple Ranking Functions." SIGIR 2025.

Rough notes/stuff we need

- Relevant links to bibliographies
- Relevant links to PISA/PyTerrier/Slack Channes/Resources/etc
- Guide on contributing to PISA
- Other codebases of interest (broader PISA/Terrier/etc projects)

Session III: Future Directions / Soapboxes

Joel's Soapbox

Long live the inverted index!

Moving towards Rust.

Sean's Soapbox

The PISA Ecosystem

CIFF

Common Index File Format CIFF is an inverted index exchange format as defined as part of the Open-Source IR Replicability Challenge (OSIRRC) initiative.

We built tools to convert:

- a CIFF blob to a PISA canonical: ciff2pisa
- a PISA canonical to a CIFF blob: pisa2ciff
- a JSONL file to a CIFF blob: jsonl2ciff

```
message Header {
 int32 version = 1:
 int32 num_postings_lists = 2;
 int32 num_docs = 3;
 int32 total_postings_lists = 4;
 int32 total_docs = 5;
 int64 total_terms_in_collection = 6:
 double average_doclength = 7;
 string description = 8;
message Posting {
 int32 docid = 1:
 int32 tf = 2;
message PostingsList {
 string term = 1;
 int64 df = 2:
 int64 cf = 3:
 repeated Posting postings = 4;
message DocRecord {
 int32 docid = 1;
 string collection_docid = 2;
 int32 doclength = 3;
```

Figure 1: Protobuf definitions of messages in CIFF.

Jimmy Lin et al. Supporting Interoperability Between Open-Source Search Engines with the Common Index File Format. SIGIR 2020.

CIFF-Hub



CIFF Hub

Common Index File Format CIFF is an inverted index exchange format as defined as part of the Open-Source IR Replicability Challenge (OSIRRC) initiative.

The Ciff Hub hosts many indexes and queries for a variety of collections and models.

Python Integration

Trades off some flexibility and efficiency for convenience.

Indexing:

```
from pyterrier_pisa import PisaIndex
# index from your own data:
index = PisaIndex('my-index.pisa')
index.index([
 {'docno': '1', 'text': 'Check out the PISA engine!'},
 {'docno': '2', 'text': 'Here is the Python integration'},
# index from a huggingface dataset:
from datasets import load_dataset
index = PisaIndex('arthur-conan-doyle.pisa')
index.index(load dataset('macavaney/arthur-conan-doyle'))
```



Python Integration

Trades off some flexibility and efficiency for convenience.

Retrieval:

```
from pyterrier_pisa import PisaIndex
index = PisaIndex('msmarco-passage.pisa')
bm25 = index.bm25()
bm25.search('pisa engine')
     aid
         query
                       docno
                                 score
                                      rank
      1 pisa engine 7378605 18.529522
    1 pisa engine 49001 17.285473
   1 pisa engine 3322192 17.025841
   1 pisa engine 4800437 16.870358
      1 pisa engine 2480271 16.774212
```



Python Integration

Trades off some flexibility and efficiency for convenience.

Sharing Indexes:

```
from pyterrier_pisa import PisaIndex

# share an index to huggingface
PisaIndex.to_hf('macavaney/my-index.pisa')

# load an index from huggingface
PisaIndex.from_hf('macavaney/msmarco-passage.pisa')
```

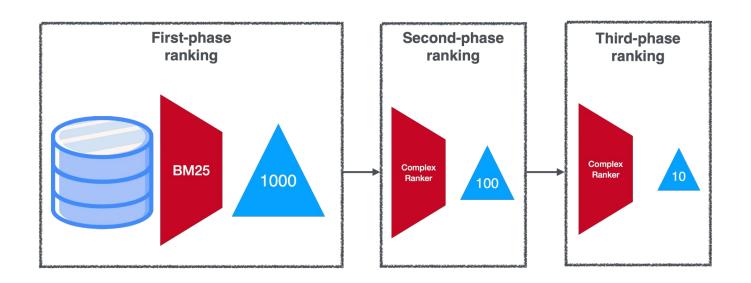


The future of the PISA engine

- We want to build a more user-friendly platform.
- Seamless notebooks one-command Colab template.
- First-class LSR support.
- MCP server lightweight micro-control-plane exposing REST/gRPC endpoints for search & index-management
- Al hooks for RAG: flexible connectors that integrate embedding, hybrid retrieval, and LLM post-processing.

Cascading Retrieval

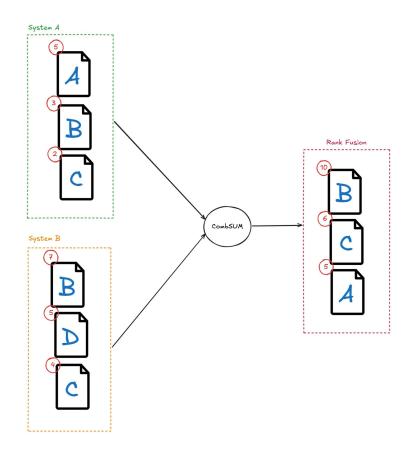
Build a multi-level architecture, from simple to complex (= cheap to expensive)



Hybrid Retrieval

Perform hybrid retrieval via rank fusion mechanisms

Fusion for Information Retrieval is the the process of combining multiple sources of information to produce a single result list in response to a query.



Reranking

Two-stage retrieval process inverted index-based candidate generation and multi-vector reranking.

Table 3. Effectiveness metrics and mean response time (MRT, in ms) for top-10 retrieval using PLAID vs. two-stage on Dev Queries, TREC 2019, and TREC 2020.

-	Dev		TREC 20	019	TREC 2020		
	MRR	MRT	nDCG@10	MRT	nDCG@10	MRT	
ColBERT	39.99	51.25	74.26	51.46	73.99	50.21	
ESPLADE	38.75	3.07	71.33	3.13	71.14	3.20	
$+ \ ConstBERT_{32}$	39.52	4.95	74.38	5.50	74.33	5.23	